COUNTING SOLUTIONS TO BINOMIAL COMPLETE INTERSECTIONS

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Abstract. We show that the problem of counting the total number of affine solutions (with and without multiplicities) of a system of n binomials in n variables is $\#P$-hard. We use commutative algebra tools to reduce the study of these solutions to a combinatorial problem on a graph associated to the exponents occurring in the given binomials.

1. Introduction

A *binomial ideal* in the polynomial ring $k[x_1, \ldots, x_n]$ over a field k, is an ideal generated by binomials: $ax^\alpha - bx^\beta$, where $\alpha, \beta \in \mathbb{N}^n$ and $a, b \in k^*$. Binomial ideals are quite ubiquitous in very different contexts particularly those involving toric geometry and its applications [9, 28], in the study of semigroup algebras, and in the modern versions of hypergeometric systems of differential equations [25, 6]. While binomial ideals are quite amenable to Gröbner and standard bases techniques [19, 20], they also provide some of the "worst-case" examples in computational algebra, such as the Mayr-Meyer ideals [22].

In this paper we consider ideals generated by n binomials in $R := k[x_1, \ldots, x_n]$, with $\mathrm{char}(k) = 0$. We are interested in determining when they define a complete intersection and, in this case, to count the number of solutions over the algebraic closure $\bar{k}$ of k. We will obtain properties of these ideals directly in terms of the given data: the exponents α, β, and the coefficients a, b.

Our main object of study is a system of binomials with non-zero coefficients. Thus, we may assume that they are of the form

$$p_j(c; x) \ := \ x^{\alpha_j} - c_j x^{\beta_j}\,; \quad j = 1, \ldots, n, \tag{1.1}$$

where $\alpha_j, \beta_j \in \mathbb{N}^n$, $\alpha_j \neq \beta_j$. Let $\mathcal{J}$ be the ideal generated by $p_1, \ldots, p_n$ in the polynomial ring $k(c)[x]$. Given a choice of coefficients $c \in (k^*)^n$, let $\mathcal{J}_c$ be the ideal in R generated by $p_1(c; x), \ldots, p_n(c, x)$ and $\mathbb{V}_c \subset \bar{k}^n$ the variety defined by $\mathcal{J}_c$.

We will say that $p_1, \ldots, p_n$ is a generic complete intersection if $\mathcal{J}_c$ is a complete intersection in the polynomial ring R, i.e. $\mathbb{V}_c$ is a finite set, for generic coefficients $c \in (k^*)^n$. We will abbreviate this by saying that $p_1, \ldots, p_n$ is a *gci*. We will describe in Section 2 how to decide whether $p_1, \ldots, p_n$ is a gci directly from the exponents α_j, β_j and how to do this in polynomial time.

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If $\mathcal{J}_c$ is a complete intersection for some $c \in (k^*)^n$, then $\mathbb{V}_c \cap (\bar{k}^*)^n \neq \emptyset$ if and only if the $n \times n$ matrix

$$B \;:=\; \begin{pmatrix} \alpha_1 - \beta_1 \\ \alpha_2 - \beta_2 \\ \vdots \\ \alpha_n - \beta_n \end{pmatrix}, \tag{1.2}$$

whose j-th row is the vector $\alpha_j - \beta_j$, is non-singular. In this case, $\mathcal{J}_c$ is a complete intersection for *all* choices of coefficients $c \in (k^*)^n$ and the number of points in $\mathbb{V}_c \cap (\bar{k}^*)^n$, counted with multiplicity, equals $|\det B|$ by Bernstein's Theorem. Even if computing the number of zeros in the torus is easy and bounding the total number of zeros is equally straightforward by Bézout's Theorem, it turns out that counting the number of solutions, with or without multiplicity, is a $\#P$-hard problem. Indeed, we show that particular instances correspond to counting independent sets in bipartite graphs, or more generally, antichains in a poset, which are known to be $\#P$-complete problems [31, 24]. The number of solutions in $\bar{k}^n$, counted with multiplicity, is given by the associated *stable mixed volume* [17]. Hence the computation of the stable mixed volume of binomials is a $\#P$-hard problem as well. In Section 4 we identify some cases where the counting problem can be solved in polynomial time.

Given a generic complete intersection $p_1, \ldots, p_n$, let

$$d \;:=\; \dim_k k[x_1, \ldots, x_n]/\mathcal{J}_c \;; \quad D \;:=\; \dim_k k[x_1, \ldots, x_n]/\sqrt{\mathcal{J}_c} \tag{1.3}$$

be the total number of points in the variety $\mathbb{V}_c$, counted with and without multiplicity. Given an index set $L \subset \{1, \ldots, n\}$, we denote by μ_L, the number of points in $\mathbb{V}(\mathcal{J}) \cap \bar{k}^n_L$, $\bar{k}^n_L \;:=\; \{x \in \bar{k}^n : x_\ell = 0 \text{ if and only if } \ell \in L\}$, counted with multiplicity. We set $[n] := \{1, \ldots, n\}$ and $\mu := \mu_{[n]}$, the multiplicity at the origin.

In Section 3 we compute d, D, and μ_L for a gci. A key ingredient is what we call *parametric reduction* which allows us to reduce the study of generic complete intersection binomial ideals to a particular class of ideals with a normalized presentation. We show in Theorem 3.2 that we can keep track of the various multiplicities through the process of parametric reduction. We then compute d and D for so-called *irreducible systems*. We show that an irreducible system which is in normal form may behave in one of three possible ways: its binomials are a standard basis for either a global or a local term order, or they are weighted homogeneous. This allows us to read off the dimension and multiplicities from the exponents (cf. Theorem 3.5). Interestingly, the linear algebra problem that underlies these results appeared in the work of Vinberg about Cartan matrices [18, Theorem 4.3]. For generic exponents, a binomial system in normal form is irreducible and has $\det B \neq 0$. Hence, Theorem 3.5 may be viewed as furnishing a polynomial time algorithm for computing the number of solutions of a complete intersection binomial system with generic exponents and arbitrary non-zero coefficients.

We complete Section 3 by considering the case of a general gci. Using a well-known quadratic-time algorithm, due to Tarjan [30], we find a block decomposition of the system into irreducible ones. This decomposition leads to an acyclic directed graph and we show in Theorem 3.15 how the dimensions and multiplicities may be computed from this graph.

The last section is devoted to the counting complexity issues already mentioned. Here we reverse the correspondence from binomial systems to acyclic digraphs and assign to each such graph a simple binomial system whose number of solutions correspond to invariants of the graph whose computation is known to be $\#P$-complete.

We end this introduction with a brief discussion of some of the problems which motivated our study. An important subfamily of binomial ideals is given by the toric ideals associated to configurations $A = \{a_1, \ldots, a_m\} \subset \mathbb{Z}^k$ of integral points spanning $\mathbb{Z}^k$:

$$I_A = \langle x^u - x^v \ ; \ A \cdot (u - v) = 0 \rangle ,$$

where $u, v \in \mathbb{N}^m$. In particular, beginning with the work of Herzog [15] and Delorme [5] the question of classifying complete intersection toric ideals (and the corresponding semigroup algebras) has been extensively studied by many authors [1, 4, 10, 11, 12, 26]. A key step in many of these works is the study of the ideal generated by binomials $x^{u_i} - x^{v_i}$ associated with a $\mathbb{Z}$-basis of the kernel of A. More generally, given $\mathbb{Q}$-linearly independent elements $\nu_1, \ldots, \nu_r \in \mathbb{Z}^m$, consider the associated lattice basis ideal $J \subset k[x_1, \ldots, x_m]$, generated by the binomials

$$b_j = x^{u_j} - x^{v_j} \ ; \quad j = 1, \ldots, r,$$

where $\nu_j = u_j - v_j$, and $u_j, v_j \in \mathbb{N}^m$ have disjoint support. Let $\mathcal{L} \subset \mathbb{Z}^m$ denote the lattice spanned by $\nu_1, \ldots, \nu_r$ and let $I_{\mathcal{L}} := \langle x^u - x^v : u - v \in \mathcal{L} \rangle$ be the corresponding lattice ideal. We assume that these ideals are homogeneous, i.e. $w_1 + \cdots + w_m = 0$, for every $w \in \mathcal{L}$.

The ideal $I_{\mathcal{L}}$ is prime if and only if the lattice $\mathcal{L}$ is saturated. If $\mathcal{L}$ is not saturated, then $I_{\mathcal{L}}$ has g radical primary components, where g is the index of $\mathcal{L}$ in its saturation. Moreover, all these components have the same degree, equal to the degree $d_{\mathcal{L}}$ of the associated toric variety [9].

We can apply Theorem 3.15 to compute the multiplicity and geometric degree [2] of the primary components of J. This may be used to describe the holonomic rank of Horn systems of hypergeometric partial differential equations and to study sparse discriminants, generalizing the codimension-two case. [7, 6].

A straightforward extension of the results of [16] to non-saturated lattices gives the following description of all primary components $\mathfrak{q}$ of J. Let $K \subset \{1, \ldots, m\}$ and $Z(K) \subset \{1, \ldots, r\}$ as in (2.3). Assume that $n := |Z(K)| = |K|$ and for all $j \notin Z(K)$

$$\operatorname{supp}(u_j) \cap K = \operatorname{supp}(v_j) \cap K = \emptyset.$$

Let $\mathfrak{p}'$ be a primary component of the lattice ideal $I_{\mathcal{L}'}$ associated to the sublattice of $\mathbb{Z}^{m-n}$ spanned by ν_j, $j \notin Z(K)$. Then, the ideal

$$\mathfrak{q} \ = \ \mathfrak{p}' \ + \ \langle b_i, i \in Z(K) \rangle$$

is a primary component of J with associated prime

$$\mathfrak{p} \ = \ \mathfrak{p}' \ + \ \langle x_k, k \in K \rangle.$$

Note that for $K = \emptyset$ we recover the components of $I_{\mathcal{L}}$.

In order to describe the multiplicity and geometric degree of a component $\mathfrak{q}$, let us assume that $K = Z(K) = \{1, \ldots, n\}$ and for any $w \in \mathbb{Z}^m$, denote $\pi(w) = (w_1, \ldots, w_n)$. Let $\alpha_j = \pi(u_j)$, $\beta_j = \pi(v_j)$ and set

$$p_j(c; x) \ = \ x^{\alpha_j} - c_j x^{\beta_j} \, , c_j \in k^*.$$

Since J is a complete intersection, $p_1, \dots, p_n$ is a gci. Let μ denote the multiplicity at the origin. Fix coefficients $c \in (k^*)^n$ such that $\mathcal{J}_c$ is a complete intersection. Since

$$\mu = \text{length}\,(k[x_1, \dots, x_n]/\mathcal{J}_c)_0 = \text{length}\,(k[x_1, \dots, x_m]/J)_{\mathfrak{p}}\,,$$

and the degree of $\mathfrak{p}$ equals that of $\mathfrak{p}'$, we have

Proposition 1.1. *With notation as above, the multiplicity of* $\mathfrak{q}$ *equals* μ *and the geometric degree of* $\mathfrak{q}$ *equals* $d_{\mathcal{L}'} \cdot \mu$.

As a second application, consider a system of constant coefficient partial differential equations defined by n operators of the form

$$a_j \partial^{\alpha_j} - b_j \partial^{\beta_j}\,; \quad j = 1, \dots, n, \tag{1.4}$$

where $a_j, b_j \in k^*$, $\alpha_j, \beta_j \in \mathbb{N}^n$, $\alpha_j \neq \beta_j$. Assume moreover that the ideal J in $k[x_1, \dots, x_n]$ generated by the binomials $a_j x^{\alpha_j} - b_j x^{\beta_j}$ is zero-dimensional. As before, let μ_L the number of points in $\mathbb{V}(J) \cap \bar{k}^n_L$ counted with multiplicity. From [29, Chapter 10], we have the following characterization.

Proposition 1.2. *Let* $L \subseteq \{1, \dots, n\}$*. The dimension of the space of solutions to (1.4) which depend polynomially on the variables* x_ℓ*,* $\ell \in L$*, and exponentially on the remaining variables* x_j*,* $j \notin L$*, equals* μ_L*.*

These dimensions can then be computed using the results in Section 3, particularly formula (3.14).

2. Complete Intersections and normal forms

We begin by considering the question of when binomials $p_1(c;x), \dots, p_n(c,x)$ as in (1.1) define a complete intersection when viewed as elements of the Laurent polynomial ring $S := k[x_1^{\pm 1}, \dots, x_n^{\pm 1}]$. Let B be the $n \times n$ exponent matrix defined in (1.2). We note that even though the rows of B are only defined up to sign, this will not affect our arguments. It follows from [9, Theorem 2.1] that if $\det B \neq 0$ then, for any choice of coefficients in $(k^*)^n$, $p_1(c;x), \dots, p_n(c;x)$ define a regular sequence in S. Moreover, the system of equations

$$p_j(c;x) \;=\; 0\,; \quad j = 1, \dots, n \tag{2.1}$$

has $|\det B|$-many solutions in the algebraic torus $(\bar{k}^*)^n$ and all of them are simple.

On the other hand, if $\det B = 0$ then $p_1(c;x), \dots, p_n(c;x)$ does not define a complete intersection in S for any choice of coefficients. Indeed, if the system (2.1) has a solution $x \in (\bar{k}^*)^n$, it will necessarily have infinitely many. Let $\mathcal{R}$ be the lattice of relations

$$\mathcal{R} \;:=\; \{m \in \mathbb{Z}^n : \sum_{j=1}^n m_j(\alpha_j - \beta_j) = 0\}.$$

For any $m \in \mathcal{R}$ we have a $\bar{k}^*$-action on the set of solutions of (2.1) defined by $(t;x) \mapsto (t^{m_1}x_1, \dots, t^{m_n}x_n)$, and therefore the set of solutions could never be finite. Note also that if $\det B = 0$ then, for generic coefficients c_j, (2.1) has no solutions. In fact, if $x \in (\bar{k}^*)^n$ is a solution of (2.1) we have

$$x^{\alpha_j - \beta_j} = c_j, \text{ for all } j \;=\; 1, \dots, n,$$

and therefore

$$\prod_{j=1}^{n} c_j^{m_j} \;=\; 1\,, \text{ for all } m \in \mathcal{R}.$$

Thus, if $\nu^1, \dots, \nu^r$ is a basis of $\mathcal{R}$, a necessary condition for $p_1(c;x), \dots, p_n(c,x)$ to have a solution in $(\bar{k}^*)^n$ is that,

$$\prod_{j=1}^{n} c_j^{\nu_j^\ell} \;=\; 1 \;\text{ for all }\; \ell = 1, \dots, r. \tag{2.2}$$

This condition is also sufficient. Suppose that (2.2) holds and let $\mathcal{L}$ be the sublattice of $\mathbb{Z}^n$ spanned by $\alpha_j - \beta_j$, $j = 1, \dots, n$. Denote by $\rho\colon \mathcal{L} \to \bar{k}^*$ the group homomorphism (i.e. the partial character) defined by

$$\rho(\alpha_j - \beta_j) = c_j\,.$$

The equalities in (2.2) imply that ρ is well-defined and, since up to a monomial (which is invertible in the Laurent polynomial ring),

$$p_j(x) \;=\; x^{\alpha_j - \beta_j} - \rho(\alpha_j - \beta_j)$$

it follows from [9, Theorem 2.6] that $p_1(c;x), \dots, p_n(c,x)$ define an ideal in S of codimension equal to the rank of $\mathcal{L}$. Hence we obtain:

Proposition 2.1. *Let $p_1(c;x), \dots, p_n(c,x)$ be as in (1.1) and B as above. For any choice of coefficients $c \in (k^*)^n$, the ideal they generated in S is a complete intersection if and only if $\det B \neq 0$. If $\det B = 0$ and the identities (2.2) are satisfied then the binomials (1.1) define an ideal in S of codimension equal to the rank of B.*

In the remaining part of this section, we will discuss criteria for deciding when $p_1, \dots, p_n$ is a gci. Since we are not assuming that the supports of the vectors α_j and β_j are disjoint, the matrix B, by itself, does not allow us to recover the exponents of the binomials (1.1). It is useful to introduce the following concept, already present in the work of Scheja, Scheja, and Storch [26]:

Definition 2.2. Let $p_j = x^{\alpha_j} - c_j\, x^{\beta_j}$, $j = 1, \dots, n$, be a system of binomials in $k[x_1, \dots, x_n]$. For each index set $K \subset [n]$, let

$$Z(K) \;:=\; \{j \in [n] : \mathrm{supp}(\alpha_j) \cap \mathrm{K} \neq \emptyset \text{ and } \mathrm{supp}(\beta_j) \cap \mathrm{K} \neq \emptyset\}\,. \tag{2.3}$$

We start by showing that we can restrict ourselves to the case where $0 \in \mathbb{V}_c$. Since this property is equivalent to the statement that all exponent vectors are non zero, it is independent of the choice of coefficients. Suppose that there exist some exponent $\alpha_j = 0$ or some exponent $\beta_j = 0$. Then, all the variables appearing in x^{β_j} or x^{α_j}, respectively, are invertible modulo $\mathcal{J}$. Now, if all the variables appearing in x^{α_k} are invertible modulo $\mathcal{J}$ then so are all the variables in the corresponding monomial x^{β_k} and so on. Continuing this process we may identify all the variables invertible modulo $\mathcal{J}$ and, after reordering of variables and polynomials, we may assume that the variables $x_{r+1}, \dots, x_n$ are invertible and that the binomials $p_{s+1}, \dots, p_n$ involve only the variables $x_{r+1}, \dots, x_n$, while for $j \leq s$ both monomials x^{α_j} and x^{β_j} are divisible by at least one of the variables x_i, $i \leq r$, i.e. that $Z([r]) = [s]$. Following [12] we define:

Definition 2.3. Let $x' := (x_1, \dots, x_r)$, $c' := (c_1, \dots, c_r)$. For $j \leq s$, set

$$\hat{p}_j(c'; x') = p_j(c'; (x_1, \dots, x_r, 1, \dots, 1)) \,. \tag{2.4}$$

Then, the binomial system $\{\hat{p}_1, \dots, \hat{p}_s\} \subset k(c')[x']$ is called the *derived system* of $p_1, \dots, p_n$. We denote by $\hat{B}$ the associated $s \times r$ matrix as in (1.2).

Note that $0 \in \mathbb{V}(\hat{p}_1, \dots, \hat{p}_s)$ and that the matrix B is of the form

$$B \;=\; \begin{pmatrix} \hat{B} & * \\ 0 & B_2 \end{pmatrix}.$$

Lemma 2.4. *Assume* $p_1, \dots, p_n$ *as in (1.1) is a gci and let* r, s *be as above. Then,* $r = s$ *and* $\det(B_2) \neq 0$.

Proof. Since the variables $x_{r+1}, \dots, x_n$ are all invertible modulo $\mathcal{J}$, the system of equations $p_{s+1} = \cdots = p_n = 0$, is equivalent to the system $x^{\alpha_j - \beta_j} = c_j$, for all $j = s+1, \dots, n$. Hence, each integer relation among the vectors $\alpha_j - \beta_j$, $j = s+1, \dots, n$ imposes a polynomial condition on the coefficients as in (2.2). If $s < r$, then $n - r < n - s$ and so there exists a non trivial relation. Therefore, $p_1, \dots, p_n$ has generically no solutions, a contradiction. On the other hand, if $s > r$, or if $r = s$ and $\det(B_2) = 0$, then, generically, the system $p_{s+1}(x_{r+1}, \dots, x_n) = \cdots = p_n(x_{r+1}, \dots, n) = 0$ has either no solutions or infinitely many in $(\bar{k}^*)^{n-r}$. Since any solution of these equations may be extended to a solution of (2.1) by setting $x_1 = \cdots = x_r = 0$, we get a contradiction again. So $s = r$ and $\det(B_2) \neq 0$, as claimed. □

Proposition 2.5. *Let* $p_1, \dots, p_n$, B *be as above. Assume that* $s = r$ *and* $\det(B_2) \neq 0$. *Let* $\hat{p}_1, \dots, \hat{p}_r$ *be the derived system. Then* $p_1, \dots, p_n$ *is a gci if and only if* $\hat{p}_1, \dots, \hat{p}_r$ *is a gci.*

Proof. Assume $p_1, \dots, p_n$ is a gci and let $\mathcal{U}$ be an open dense subset of $(\bar{k}^*)^n$ such that the binomials with coefficients in $\mathcal{U}$ define a complete intersection ideal in $\bar{k}[x_1, \dots, x_n]$. It suffices to show that the intersection of $\mathcal{U}$ with the fiber $(\bar{k}^*)^r \times \{(1, \dots, 1)\}$ is also Zariski dense in the fiber. Let $a'' \in (\bar{k}^*)^{n-r}$ be such that $\mathcal{U} \cap \big((\bar{k}^*)^r \times \{a''\}\big)$ is Zariski dense. Given a common zero $\lambda'' \in (\bar{k}^*)^{n-r}$ of $p_{r+1}(a''; x), \dots, p_n(a''; x)$, the change of variables which sends x_i to itself for $i = 1, \dots, r$ and

$$x_j \mapsto x_j / \lambda_j, \quad j = r+1, \dots, n,$$

transforms the last $n - r$ polynomials into $x^{\alpha_j} - x^{\beta_j}$, $j = r+1, \dots, n$ and, for $i \leq r$, the binomial p_i into a non-zero multiple of

$$x^{\alpha_i} - (\lambda'')^{\alpha_i'' - \beta_i''}\, c_i\, x^{\beta_i},$$

where $\alpha_i'', \beta_i'' \in \mathbb{N}^{n-r}$ denote the vectors consisting of the last $n - r$ coordinates of α_i, β_i. Since this scalar transformation in the coefficient space $(\bar{k}^*)^r$ preserves Zariski dense subsets our assertion follows.

Conversely, assume that $\hat{p}_1, \dots, \hat{p}_r$ is a gci and that $\det(B_2) \neq 0$. Let φ be a non zero polynomial such that $\varphi(c') \neq 0$ for a given r-tuple of coefficients $c' = (c_1, \dots, c_r)$ implies that the corresponding polynomials $\hat{p}_1(c'; x'), \dots, \hat{p}_r(c'; x')$ define a complete intersection. Denote as before $c'' = (c_{r+1}, \dots, c_n)$ and consider the rational function

$$\psi(c', c'') = \prod_{\lambda'' \in \mathbb{V}_{c''}} \varphi((\lambda'')^{\alpha_1'' - \beta_1''}\, c_1, \dots, (\lambda'')^{\alpha_r'' - \beta_r''}\, c_r).$$

If $\psi(c', c'')$ is defined and non zero, then for any choice of the $|\det(B_2)|$-many roots λ'' of the last $n-r$ polynomials, the specialized system

$$p_1(c';(x',\lambda'')) = \cdots = p_r(c';(x',\lambda'')) = 0$$

has finitely many solutions and, consequently, $p_1, \ldots, p_n$ is a gci. $\square$

The following result is a reformulation of Theorem 2.3 in [12].

Theorem 2.6. *Let $p_1, \ldots, p_n$ be as in (1.1) and suppose that $0 \in \mathbb{V}(\mathcal{J})$. Then, $p_1, \ldots, p_n$ ia a gci if and only if $|Z(K)| \leq |K|$ for all $K \subset [n]$.*

Proof. Suppose there exists $K \subset [n]$ such that $|Z(K)| > |K|$. Assume that K is maximal with this property. After reordering, if necessary, we may assume that $K = \{r+1, \ldots, n\}$ and $Z(K) = \{s+1, \ldots, n\}$ where $s < r$. Since $0 \in \mathbb{V}(\mathcal{J})$, the maximality assumption implies that the first s binomials depend only on $x' = (x_1, \ldots, x_r)$. Thus, for a given choice of coefficients, the system

$$p_1(c;x') = \cdots = p_s(c;x') = 0 \tag{2.5}$$

is either inconsistent or its solution space has dimension at least $r-s > 0$. Since, any solution of (2.5) can be extended to a solution of the full system by setting the K-coordinates equal to zero, it follows that $p_1, \ldots, p_n$ is not a gci.

Conversely, suppose $|Z(K)| \leq |K|$ for all $K \subset [n]$. In order to show that $p_1, \ldots, p_n$ is a gci it suffices to prove that given any subset $L \subset [n]$, for generic coefficients $p_1(c;x), \ldots, p_n(c,x)$ has at most finitely many solutions with zeroes in $\bar{k}^n_L$, where

$$\bar{k}^n_L \;=\; \{x \in \bar{k}^n : x_\ell = 0 \text{ if and only if } \ell \in L\}\,. \tag{2.6}$$

Assume that for some choice of coefficients, there exists a solution in $\bar{k}^n_L$. Then, for any $i \notin Z(L)$, p_i depends only on the variables in J, the complement of L in $[n]$ and hence, since $0 \in \mathbb{V}(\mathcal{J})$, $Z(L)^c \subset Z(J)$. Since, by assumption $|Z(L)| \leq |L|$ and $|Z(J)| \leq |J|$, we deduce that

$$|L| \leq |Z(J)^c| \leq |Z(L)| \leq |L|,$$

and therefore $|Z(L)| = |L|$. Reordering we may assume that $J = Z(J) = [r]$ and let $B_1(L)$ denote the $r \times r$ exponent matrix as in (1.2). If $\det B_1(L) = 0$, then for generic coefficients the first r binomials have no solutions in $(\bar{k}^*)^r$ and hence, generically, $p_1(c;x), \ldots, p_n(c,x)$ have no solutions in $\bar{k}^n_L$. On the other hand, if $\det B_1(L) \neq 0$ then, for all choices of coefficients in $(k^*)^r$, there exists finitely many solutions of $p_1 = \cdots = p_r = 0$ in $(\bar{k}^*)^r$ and hence finitely many solutions of $p_1(c;x), \ldots, p_n(c,x)$ with zeroes exactly in L. $\square$

Remark 2.7. Note that in the proof of Theorem 2.6 we have shown that if $p_1, \ldots, p_n$ is a gci, $L \subset [n]$, and $\bar{k}^n_L$ is as in (2.6), then, for generic coefficients, there exists a solution in $\bar{k}^n_L$ if and only if $|Z(L)| = |L|$ and, after reordering so that $Z(L) = L = \{r+1, \ldots, n\}$, the binomials $p_1, \ldots, p_r$ depend only on the first r variables, and the corresponding $r \times r$ exponent matrix $B_1(L)$ is non-singular. Moreover, for generic $c \in (k^*)^n$, there are $|\det B_1(L)|$-many points (counted without multiplicity) in $\mathbb{V}_c \cap \bar{k}^n_L$. Then,

$$D \;=\; \sum_{\mu_L \neq 0} |\det B_1(L)|. \tag{2.7}$$

We will develop in Section 3 the combinatorics needed to describe all sets L with $\mu_L \neq 0$ and we shall show in Section 4 that counting the number of such sets is a $\#P$-hard problem.

Note that if and $0 \in \mathbb{V}(\mathcal{J})$, the condition that $p_1, \ldots, p_n$ is a gci depends only on the combinatorics of the exponents α_j, β_j. As remarked before, when $\det(B) \neq 0$, if $p_1, \ldots, p_n$ is a gci, then it is a complete intersection for any choice of the coefficients (as long as $c_j \in k^*$).

The variant of the Fischer-Shapiro criterion embodied in Theorem 2.6 allows us to determine whether $p_1, \ldots, p_n$ is a gci. However, this involves checking exponentially many conditions, one for each subset $K \subset [n]$. We will now show how this can be done in a number of steps which depends polynomially on n. We begin with the following simple corollary to Theorem 2.6.

Corollary 2.8. *Suppose* $p_1, \ldots, p_n$ *is a gci and* $0 \in \mathbb{V}(\mathcal{J})$. *Let*

$$\mathcal{M} \;=\; \{x^{\alpha_j}, x^{\beta_j} \;;\; j = 1, \ldots, n\}$$

denote the set of monomials appearing in $p_1, \ldots, p_n$. *Then for each* $i \in [n]$ *there exists* $r_i > 0$ *such that* $x_i^{r_i} \in \mathcal{M}$.

Proof. If for some $i \in [n]$, $x_i^{r_i} \notin \mathcal{M}$ for all $r_i > 0$, then $Z(\{1, \ldots, \hat{i}, \ldots, n\}) = [n]$, contradicting Theorem 2.6. □

One can easily give examples showing that the necessary condition in Corollary 2.8 is not sufficient to guarantee that $p_1, \ldots, p_n$ define a gci. However, the following stronger notion provides a sufficient condition.

Definition 2.9. We say that $p_1, \ldots, p_n$ are in *normal form* if and only if for all $i \in [n]$

$$p_i = x_i^{r_i} - c_i x^{\beta_i} \;; \quad r_i > 0,\, \beta_i \neq 0.$$

Note that if the system is in normal form then $0 \in \mathbb{V}(\mathcal{J})$.

Proposition 2.10. *Assume* $p_1, \ldots, p_n$ *are in normal form. Then* $p_1, \ldots, p_n$ *is a gci.*

Proof. For any $K \subset [n]$, $Z(K) \subset K$ and the result follows from Theorem 2.6. □

We will next show how to reduce ourselves to systems $p_1, \ldots, p_n$ in normal form.

2.1. **Parametric Reduction.** Let $p_1, \ldots, p_n$ be a binomial system and suppose that they satisfy the necessary condition in Corollary 2.8, but that it is not possible to relabel variables and binomials, or invert the coefficient of one or more binomials, so as to put the system in normal form. Then, after relabeling we may assume that p_n is of the form

$$p_n \;=\; x_n^{\ell} - c_n x_{n-1}^{m}\,,\;\; \ell, m > 0. \tag{2.8}$$

Let $q := \gcd(m, \ell)$ and set $m' := m/q$, $\ell' := \ell/q$. We will consider the polynomial map that sends polynomials in n variables $x_1, \ldots, x_n$ to polynomials in $n-1$ variables $u_1, \ldots, u_{n-1}$:

$$x_i \mapsto u_i,\;\; i = 1, \ldots, n-2; \quad x_{n-1} \mapsto u_{n-1}^{\ell'}; \quad x_n \mapsto u_{n-1}^{m'}. \tag{2.9}$$

Let $\tilde{p}_j$, $j = 1, \ldots, n-1$ be the image of the binomials $p_1, \ldots, p_{n-1}$. We will refer to $\tilde{p}_1, \ldots, \tilde{p}_{n-1}$ as a *parametric reduction* of $p_1, \ldots, p_n$ and denote by $\tilde{\mathcal{J}}$ the ideal they generate in $k(c_1, \ldots, c_{n-1})[u_1, \ldots, u_{n-1}]$.

Proposition 2.11. *Suppose $\tilde{p}_1, \dots, \tilde{p}_{n-1}$ is a parametric reduction of $p_1, \dots, p_n$ and let $\tilde{B}$ and B be the associated matrices. Then $|\det B| = q \cdot |\det \tilde{B}|$. Moreover, $0 \in \mathbb{V}(\mathcal{J})$ if and only if $0 \in \mathbb{V}(\tilde{\mathcal{J}})$ and, in this case, $p_1, \dots, p_n$ is a gci if and only if $\tilde{p}_1, \dots, \tilde{p}_{n-1}$ is gci.*

Proof. The matrix B is of the form

$$B = \begin{pmatrix} \tilde{b}_1 & \dots & \tilde{b}_{n-2} & \tilde{b}_{n-1} & \tilde{b}_n \\ \\ 0 & \dots & 0 & -m & \ell \end{pmatrix}$$

where $\tilde{b}_1, \dots, \tilde{b}_n$ are vectors in $\mathbb{Z}^{n-1}$. On the other hand, the matrix $\tilde{B}$ is given by

$$\tilde{B} = \begin{pmatrix} \tilde{b}_1 & \dots & \tilde{b}_{n-2} & \ell' \tilde{b}_{n-1} + m' \tilde{b}_n \end{pmatrix}$$

The first assertion now follows from a last-row expansion of $\det B$.

Suppose now that $p_1, \dots, p_n$ is not a gci. By Theorem 2.6 there exists $K \subset [n]$ such that $|Z(K)| > |K|$. If $K \subset [n-1]$, then $Z(K) \subset [n-1]$ as well and therefore by Theorem 2.6 $\tilde{p}_1, \dots, \tilde{p}_{n-1}$ is not a gci either. If $n \in K$, then taking $\tilde{K} = K \setminus \{n\}$ we get that $Z(K) \setminus \{n\} \subset Z(\tilde{K})$. Hence $|Z(\tilde{K})| > |\tilde{K}|$ and $\tilde{p}_1, \dots, \tilde{p}_{n-1}$ is not a gci.

Conversely, if $\tilde{p}_1, \dots, \tilde{p}_{n-1}$ is not a gci then there exists $\tilde{K} \subset [n-1]$ such that $|Z(\tilde{K})| > |\tilde{K}|$. If $\tilde{K} \subset [n-2]$ we take $K = \tilde{K}$ and then $Z(K) = Z(\tilde{K})$; if, on the other hand, $n-1 \in \tilde{K}$, then we take $K = \tilde{K} \cup \{n\}$ in which case $Z(K) = Z(\tilde{K}) \cup \{n\}$. In either case $|Z(K)| > |K|$ and we are done. □

The results of this section may be summarized in a polynomial-time algorithm to check whether a binomial system is a gci.

Theorem 2.12. *We may decide in polynomial time whether $p_1, \dots, p_n$ is a gci. Moreover, if it is known that $\det B \neq 0$ we can check if $p_1(c; x), \dots, p_n(c, x)$ is a complete intersection in time $O(n^2)$.*

Proof. It is easy to see from the procedure for constructing the derived system that this step may be accomplished in at most $O(n^2)$ steps. If the number of non-invertible variables does not equal the number of binomials in the derived system then, by Proposition 2.5, $p_1, \dots, p_n$ is not a gci. By Lemma 2.4 we next check whether $\det B_2 \neq 0$ (this is, of course, unnecessary if it is known that $\det B \neq 0$). If so, we search for binomials of the form $x_i^{r_i} - c x_j^{r_j}$ and we do parametric reduction. This process stops after a quadratic number of steps. Then $p_1, \dots, p_n$ is a gci if and only if the resulting system may be put in normal form trough reordering and relabeling; this step again requires at most quadratically-many steps. □

Example 2.13. Consider the following binomials in $k[x_1, \dots, x_8]$:

$$p_1 = x_1^2 - x_2^3; \quad p_2 = x_1 x_2 - x_1 x_3; \quad p_3 = x_1^2 x_2 x_3 - x_3^7;$$
$$p_4 = x_4^2 - x_1^2 x_4^3; \quad p_5 = x_5^2 - x_6^4; \quad p_6 = x_5 x_6 - x_2 x_3 x_7^2 x_8;$$
$$p_7 = x_5 x_7 - x_7^2; \quad p_8 = x_8^3 - x_1 x_6 x_7 x_8,$$

where, since $\det B \neq 0$, we have set all coefficients $c_j = 1$. Although the system satisfies the necessary condition in Corollary 2.8, it is not in normal form. We may apply parametric reduction simultaneously to the binomials p_1 and p_5 by considering the polynomial map from $k[x_1,\dots,x_8]$ to $k[u_1,\dots,u_6]$ that sends:

$$x_1 \mapsto u_1^3; \quad x_2 \mapsto u_1^2; \quad x_3 \mapsto u_2; \quad x_4 \mapsto u_3;$$
$$x_5 \mapsto u_4^2; \quad x_6 \mapsto u_4; \quad x_7 \mapsto u_5; \quad x_8 \mapsto u_6.$$

Here we have taken into account that the gcd of the exponents in p_5 is 2. After changing signs when necessary, the new system $\tilde{p}_1,\dots,\tilde{p}_6$ is in normal form:

$$\tilde{p}_1 = u_1^5 - u_1^3u_2; \quad \tilde{p}_2 = u_2^7 - u_1^8u_2;$$
$$\tilde{p}_3 = u_3^2 - u_1^6u_3^3; \quad \tilde{p}_4 = u_4^3 - u_1^2u_2u_5^2u_6;$$
$$\tilde{p}_5 = u_5^2 - u_4^2u_5; \quad \tilde{p}_6 = u_6^3 - u_1^3u_4u_5u_6.$$

Thus, we conclude that $p_1,\dots,p_8$ defines a complete intersection. We will compute the numerical invariants of this system in Example 3.17.

3. Computing the number of solutions

We recall that if $p_1,\dots,p_n$ is a gci then we denote by d (respectively D) the number of points in $\mathbb{V}_c \cap \bar{k}^n$ counted with multiplicity (respectively without multiplicity), for a generic choice of non-zero coefficients. Similarly, recall that for any index set $L \subset \{1,\dots,n\}$ we denote by μ_L the number of points in $\mathbb{V}_c \cap \bar{k}^n_L$ counted with multiplicity, where $\bar{k}^n_L$ is the set of points in affine space whose coordinates $x_\ell, \ell \in L$, vanish. In particular, $\mu = \mu_{[n]}$ denotes the multiplicity at the origin.

If $p_1,\dots,p_n$ is a gci but $0 \notin \mathbb{V}(\mathcal{J})$, then it follows from Lemma 2.4 and Proposition 2.5 that the invariants of $p_1,\dots,p_n$ are obtained from those of the derived system by multiplying times $|\det B_2|$. We will assume from now on that no variable is invertible modulo $\mathcal{J}$, i.e. that $0 \in \mathbb{V}(\mathcal{J})$.

We begin this section by showing that it is enough to compute the desired numerical invariants d, D, μ_L, for ideals in normal form. We then show that if the system is irreducible, in a sense made precise below, then the only zero outside the torus is the origin and its multiplicity may be easily computed from the exponents of the system. Finally, we consider the general case and show how the various dimensions depend on the combinatorics of the irreducible components.

3.1. **Multiplicities and parametric reduction.** Suppose $p_1(c;x),\dots,p_n(c,x)$ is as in (1.1) with $p_n = x_n^\ell - c_n x_{n-1}^m$, $\ell, m > 0$. Let $q = \gcd(\ell, m)$ and

$$p'_n \;=\; x_n^{\ell'} - c'_n x_{n-1}^{m'}\,.$$

We will denote by d', D', μ'_L the corresponding invariants for $p_1,\dots,p_{n-1},p'_n$.

We show, first of all, that by keeping track of q we may assume without loss of generality that m and ℓ are coprime.

Lemma 3.1. *With notation as above, set $m' = m/q$, $\ell' = \ell/q$, $p'_n = x_n^{\ell'} - c'_n x_{n-1}^{m'}$ and let B and B' be the corresponding matrices.*

1. $|\det B| \;=\; q \cdot |\det B'|$.
2. $p_1,\dots,p_n$ *is a gci if and only if* $p_1,\dots,p_{n-1},p'_n$ *is a gci.*
3. *For any index set* $L \subset \{1,\dots,n\}$, $\mu_L \;=\; q\cdot \mu'_L$.
4. $d = q\cdot d'$ *and* $D = q \cdot D'$.

Proof. The first assertion is trival while the second one follows from Theorem 2.6. In order to prove assertion 3, let $(c_1, \ldots, c_n) \in (k^*)^n$ be such that $\mathcal{J}_c$ is a complete intersection and decompose

$$p_n = x_n^\ell - c_n x_{n-1}^m = \prod_{\xi \in W_q} (x_n^{\ell'} - \xi\, x_{n-1}^{m'}) \,, \tag{3.1}$$

where W_q denotes the q-th roots of c_n. For any $\lambda \in \mathbb{V}_c$, we have

$$\dim_k (R_\lambda/(\mathcal{J}_c)_\lambda) \;=\; \sum_{\xi \in W_q} \dim_k (R_\lambda/(\mathcal{J}_\xi)_\lambda) \;,$$

where $\mathcal{J}_\xi := \langle p_1(c;x), \ldots, p_{n-1}(c;x), x_n^{\ell'} - \xi\, x_{n-1}^{m'} \rangle$. Therefore,

$$\sum_{\lambda \in \mathbb{V}_c \cap \bar{k}_L^n} \dim_k (R_\lambda/(\mathcal{J}_c)_\lambda) \;=\; \sum_{\xi \in W_q} \;\sum_{\lambda \in \mathbb{V}(\mathcal{J}_\xi) \cap \bar{k}_L^n} \dim_k (R_\lambda/(\mathcal{J}_\xi)_\lambda) \;,$$

By a scalar change of variables it follows that

$$\sum_{\lambda \in \mathbb{V}(\mathcal{J}_\xi) \cap \bar{k}_L^n} \dim_k (R_\lambda/(\mathcal{J}_\xi)_\lambda) \;,$$

is independent of $\xi \in W_q$ and, since it agrees with μ'_L, we obtain that

$$\mu_L \;=\; q\, \mu'_L$$

as claimed. The last assertion follows directly from the previous one and the factorization (3.1). ☐

We next show that multiplicities are not altered under parametric reduction. If the binomial system $p_1, \ldots, p_n$ is a gci, and $p_n = x_n^\ell - c_n x_{n-1}^m$, $\ell, m > 0$ coprime, let $\tilde{p}_1, \ldots, \tilde{p}_{n-1}$ be the binomial system obtained through parametric reduction. We will denote by $\tilde{d}, \tilde{D}$ and $\mu_{\tilde{L}}$ the corresponding invariants.

Given $L \subset [n]$ we denote by $\tilde{L} := L \cap [n-1]$. Conversely, given $\tilde{L} \subset [n-1]$ set $L = \tilde{L}$ if $n-1 \notin \tilde{L}$ and $L = \tilde{L} \cup \{n\}$ otherwise. Note that the correspondence $L \mapsto \tilde{L}$ establishes a bijection between index sets $L \subset [n]$ such that $\mu_L \neq 0$ and subsets $\tilde{L} \subset [n-1]$ such that $\mu_{\tilde{L}} \neq 0$.

Theorem 3.2. *Suppose that $p_1, \ldots, p_n$ is a gci and $p_n = x_n^\ell - c_n x_{n-1}^m$, with ℓ, m coprime positive integers. Let $\tilde{p}_1, \ldots, \tilde{p}_{n-1}$ be the binomial system obtained through parametric reduction. Then, for any $L \subset [n]$*

$$\mu_L = \mu_{\tilde{L}} \tag{3.2}$$

and, consequently, $d = \tilde{d}$. Moreover, $D = \tilde{D}$.

Proof. Let $c = (c_1, \ldots, c_n) \in (k^*)^n$ be such that $\mathcal{J}_c$ is a complete intersection. We may assume without loss of generality that $c_n = 1$. Let $\tilde{c} = (c_1, \ldots, c_{n-1})$ and denote by $\tilde{\mathcal{J}}_{\tilde{c}}$ the ideal generated by $\tilde{p}_1(\tilde{c};u), \ldots, \tilde{p}_{n-1}(\tilde{c};u)$ in the ring $k[u]$. Given any $\tilde{\lambda} = (\lambda_1, \ldots, \lambda_{n-1}) \in \mathbb{V}(\tilde{\mathcal{J}}_{\tilde{c}}) \subset \bar{k}^{n-1}$, let us denote by λ the point $(\lambda_1, \ldots, \lambda_{n-2}, \lambda_{n-1}^\ell, \lambda_{n-1}^m) \in \mathbb{V}(\mathcal{J}_c)$. It suffices to prove that at the level of local rings

$$\dim_{\bar{k}} (R \otimes_k \bar{k})_\lambda/(\mathcal{J}_c)_\lambda \;=\; \dim_{\bar{k}} (\tilde{R} \otimes_k \bar{k})_{\tilde{\lambda}}/(\tilde{\mathcal{J}}_{\tilde{c}})_{\tilde{\lambda}}. \tag{3.3}$$

We will denote by A_1 the localization of $\bar{k}[u_1, \ldots, u_{n-1}]$ at $\tilde{\lambda}$ and by A_2 the localization of $\bar{k}[u_1, \ldots, u_{n-2}, u_{n-1}^\ell, u_{n-1}^m]$ at $\tilde{\lambda}$. Let $(\hat{\mathcal{J}}_{\tilde{c}})_{\tilde{\lambda}}$ be the ideal generated by

$\tilde{p}_1(\tilde{c};u),\ldots,\tilde{p}_{n-1}(\tilde{c};u)$ in A_2 so that $(\tilde{\mathcal{J}}_{\tilde{c}})_{\tilde{\lambda}} = A_1\cdot(\hat{\mathcal{J}}_{\tilde{c}})_{\tilde{\lambda}}$. Since m and ℓ are coprime it is clear that

$$\dim_{\bar{k}}(R\otimes_k \bar{k})_\lambda/(\mathcal{J}_c)_\lambda \;=\; \dim_{\bar{k}} A_2/(\hat{\mathcal{J}}_{\tilde{c}})_{\tilde{\lambda}}\,.$$

Thus, the result will follow if we show that

$$\dim_{\bar{k}} A_1/(\tilde{\mathcal{J}}_{\tilde{c}})_{\tilde{\lambda}} \;=\; \dim_{\bar{k}} A_2/(\hat{\mathcal{J}}_{\tilde{c}})_{\tilde{\lambda}} \tag{3.4}$$

The following proof of (3.4) was suggested to us by Mircea Mustata.

We recall from [21, §14] the following notion of *multiplicity*: Let $(R,\mathfrak{m})$ be a d-dimensional Noetherian local ring, M a finite R-module and $\mathfrak{q}$ an $\mathfrak{m}$-primary ideal. The multiplicity of M with respect to $\mathfrak{q}$ equals

$$\mathfrak{e}(\mathfrak{q},M) \;=\; \lim_{m\to\infty}\frac{d!}{m^d}\,\mathrm{length}(M/\mathfrak{q}^{m+1}M) \tag{3.5}$$

Since both A_1 and A_2 are Cohen-Macaulay rings of dimension $n-1$ and

$$\tilde{p}_1(\tilde{c};u),\ldots,\tilde{p}_{n-1}(\tilde{c};u)$$

define a regular sequence in A_2, hence in A_1 as well, it follows from [21, Theorem 14.11] that

$$\dim_{\bar{k}} A_2/(\hat{\mathcal{J}}_{\tilde{c}})_{\tilde{\lambda}} \;=\; \mathfrak{e}((\hat{\mathcal{J}}_{\tilde{c}})_{\tilde{\lambda}},A_2) \quad\text{and}\quad \dim_{\bar{k}} A_1/(\tilde{\mathcal{J}}_{\tilde{c}})_{\tilde{\lambda}} \;=\; \mathfrak{e}((\tilde{\mathcal{J}}_{\tilde{c}})_{\tilde{\lambda}},A_1).$$

On the other hand, A_1 may be considered as a A_2-module and it is clear from (3.5) that

$$\mathfrak{e}((\tilde{\mathcal{J}}_{\tilde{c}})_{\tilde{\lambda}},A_1) \;=\; \mathfrak{e}((\hat{\mathcal{J}}_{\tilde{c}})_{\tilde{\lambda}},A_1)\,.$$

Finally, [21, Theorem 14.8] gives that

$$\mathfrak{e}((\hat{\mathcal{J}}_{\tilde{c}})_{\tilde{\lambda}},A_1) = \mathrm{rank}_{A_2}A_1\cdot\mathfrak{e}((\hat{\mathcal{J}}_{\tilde{c}})_{\tilde{\lambda}},A_2) = \mathfrak{e}((\hat{\mathcal{J}}_{\tilde{c}})_{\tilde{\lambda}},A_2),$$

since the assumption that m and ℓ are coprime implies that the two domains A_1, A_2 have the same fraction field and so $\mathrm{rank}_{A_2}A_1 = 1$. This proves (3.4). ☐

3.2. **Irreducible Systems.**

Definition 3.3. A binomial system $p_1,\ldots,p_n$ is said to be irreducible if it is in normal form and it is not possible to reorder it so as to find a proper index subset $I\subset[n]$ such that for every $i\in I$ the binomial p_i depends only on the variables $x_j, j\in I$.

Recalling that a system in normal form is a gci and that $0\in\mathbb{V}(\mathcal{J})$, we easily have:

Lemma 3.4. *Let $p_1,\ldots,p_n$ be an irreducible system as in (1.1) and let $c\in(k^*)^n$ be such that $\mathcal{J}_c$ is a complete intersection. Then if $a\in\mathbb{V}_c$, either $a=0$ or $a\in(\bar{k}^*)^n$.*

Proof. Given $a\in\mathbb{V}_c$, let $I=\{i\in[n]: a_i\neq 0\}$. If $i\in I$ then, since $p_1,\ldots,p_n$ is in normal form,

$$p_i(c;x) = x_i^{r_i} - c_i x^{\beta_i}\,;\; r_i>0,\;\beta_i\neq 0,$$

and, since $a_i\neq 0$, we must have $\mathrm{supp}(\beta_i)\subset I$ for all $i\in I$. This contradicts the irreducibility of $p_1,\ldots,p_n$ unless $I=[n]$ or $I=\emptyset$. ☐

The following theorem identifies d and μ for irreducible systems. Our arguments are built on the proof of a result of Vinberg (cf. [18, Theorem 4.3]).

Theorem 3.5. *Given an irreducible system*

$$p_i(c;x) \ = \ x_i^{r_i} - c_i x^{\beta_i}\ ,\ i = 1, \ldots, n,$$

where $r_i > 0,\ \beta_i \in \mathbb{N}^n,\ \beta_i \neq 0$, *then:*

- *If all principal minors of* B *are positive*

  $$d = r_1 \cdots r_n\ ; \quad \mu = d - |\det B|.$$

  Such a system will be called a global *irreducible system.*
- *Otherwise,* $\mu = r_1 \cdots r_n$ *and* $d = \mu + |\det B|$. *In this case we say that the system is* local.

Proof. Let us fix throughout coefficients $c \in (k^*)^n$ such that $\mathcal{J}_c$ is a complete intersection. Since the system is in normal form, the entries of B are $b_{ii} = r_i - (\beta_i)_i$ and $b_{ij} = -(\beta_i)_j, i \neq j$. Hence, its off-diagonal terms are non-positive. Moreover, the irreducibility of the system implies that B is indecomposable in the sense of [18]. In fact, the irreducibility of the system implies a stronger condition, namely [18, Lemma 4.3]: Suppose $u \in \mathbb{R}^n$ is a vector with non-negative entries and that $B \cdot u \geq 0$ in the sense that all its entries are non-negative as well. Then either $u = 0$, or $u > 0$, i.e. all its entries are strictly positive. Indeed, let $I = \{i \in [n] : u_i = 0$, then for any $i \in I$, $(B \cdot u)_i \leq 0$ and equality occurs if and only if $b_{ij} = 0$ for all $j \notin I$. Hence, by irreducibility we must have $I = [n]$ or $I = \emptyset$.

Given that [18, Lemma 4.3] holds in our case, we can apply Theorem 4.3 in [18] and conclude that three cases are possible:

- There exists $w \in \mathbb{Q}^n$ all of whose entries are positive such that $B \cdot w > 0$.
- There exists $w \in \mathbb{Q}^n$, all of whose entries are positive such that $B \cdot w < 0$.
- $\operatorname{rank}(B) = n - 1$ and there exists $w \in \mathbb{Q}^n$ all of whose entries are positive such that $B \cdot w = 0$.

According to [3, Theorem 2.3], the first condition is equivalent to the statement that all principal minors of B are positive which implies, in particular, that all the diagonal entries of B are strictly positive. These are the so-called M-matrices of [3]. Moreover, if we consider a term order in $k[x_1, \ldots, x_n]$ which refines the weight order defined by w, the term $x_i^{r_i}$ will be the leading term in $p_i(c;x)$, and hence $p_1(c;x), \ldots, p_n(c,x)$ is a Gröbner basis. It then follows that $d = r_1 \cdots r_n$ and, by Lemma 3.4, $\mu = d - |\det B|$.

In the second case we can similarly define a local order (cf. [13]) for which the leading term of $p_i(c;x)$ is $x_i^{r_i}$. Hence $p_1(c;x), \ldots, p_n(c,x)$ is a standard basis in the local quotient ring at the origin and, consequently, $\mu = r_1 \cdots r_n$ and $d = \mu + |\det B|$. We note that this is valid wether $\det B \neq 0$ or $\det B = 0$ since, in the latter case, $\mathcal{J}_c$ is a complete intersection if and only if $\mathbb{V}_c = \{0\}$.

In the third case, the binomials $p_i(c;x)$ are weighted homogeneous relative to the weight w and therefore $\mu = r_1 \cdots r_n$ and $d = \mu + \det B$ since, again, $\mathbb{V}_c$ consists of only the origin. Thus this case behaves as the previous one and we will also refer to it as a local case. □

Remark 3.6. We note that if $n = 1$, the system $p = x^\alpha - cx^\beta, \alpha \neq \beta$, will be local if $\alpha < \beta$ and global if $\alpha > \beta$.

3.3. **The General Case.** We consider now general gci systems in normal form. Throughout this subsection we will, again, fix coefficients $c \in (k^*)^n$ so that $\mathcal{J}_c$ is a complete intersection. For economy of notation we will denote simply by p_i the corresponding binomials in $k[x_1, \ldots, x_n]$. If the system $p_1, \ldots, p_n$ is not irreducible, then, as Lemma 3.8 shows, it is possible to choose an increasing sequence

$$0 = \nu_0 < \nu_1 < \cdots < \nu_s = n \tag{3.6}$$

so that if $I_a = \{\nu_{a-1} + 1, \ldots, \nu_a\}$, then the following holds:

- For $i \in I_a$, $p_i \in k[x_j; j \in I_1 \cup \cdots \cup I_a]$.
- The system $\hat{p}_i := p_i(1, \ldots, 1, x_{\nu_{a-1}+1}, \ldots, x_{\nu_a})$, $i \in I_a$, is irreducible.

Definition 3.7. A system of this form will be said to be in *triangular* form relative to the blocks $I_1, \ldots, I_s$. Given a reducible system in triangular form, we will refer to the system $\{\hat{p}_i, i \in I_a\}$ as the restriction of $p_1, \ldots, p_n$ to I_a and denote it, for short, by $\hat{p}^a$.

Lemma 3.8. *Any system of n binomials $p_1, \ldots, p_n$ in normal form (2.9) can be put in triangular form in time $O(n^2)$.*

Proof. Consider the ocurrence matrix N: this is a 0-1 matrix with $n_{ij} \neq 0$ if and only if $i \neq j$ and p_i depends on x_j (i.e. if $p_i = x_i^{r_i} - c_i x^{\beta_i}$ with $\beta_{ij} \neq 0$). This is a standard construction, first used by Steward [27], for the analysis of the structure of large systems of equations. Note that, because the system is in normal form, putting $p_1, \ldots, p_n$ in triangular form corresponds precisely to finding a permutation matrix P such that ${}^tPNP$ is block lower triangular, with the irreducible subsystems of $p_1, \ldots, p_n$ corresponding to the irreducible diagonal square blocks along the diagonal of ${}^tPNP$.

Tarjan's algorithm [30] to search for the strongly connected components of the directed graph associated to N provides an efficient method for finding such permutation matrix P [8, 23]; it runs in time linear in the number of vertices plus the number of edges of the graph. $\square$

Given a system in normal form and triangular relative to $I_1, \ldots, I_s$, let $\delta_a = |\det B_a|$, where B_a is the matrix associated with the system $\hat{p}^a$ and

$$\rho_a = \prod_{j \in I_a} r_j .$$

We also denote by μ_a the multiplicity of $\hat{p}^a$ at 0 and by d_a the total number of solutions of $\hat{p}^a$ counted with multiplicity.

For a triangular system $p_1, \ldots, p_n$, its associated matrix is block lower-triangular:

$$B = \begin{pmatrix} B_1 & 0 & \ldots & 0 \\ C_{21} & B_2 & \ldots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ C_{s1} & C_{s2} & \ldots & B_s \end{pmatrix} . \tag{3.7}$$

The number of solutions of the system $p_1, \ldots, p_n$ and the patterns of possible zero coordinates of the solutions are best described in terms of the directed acyclic graph G with s vertices labeled $\{1, \ldots, s\}$ and an arrow from node a to node b if and only if the rectangular submatrix C_{ba} is not identically zero. We recall that a vertex is called a *source* if it is not the head of any arrow. The subset of sources of the vertex set of a subgraph H of G will be denoted by $S(H)$.

Remark 3.9. We can think of G as a weighted graph, where each vertex $a \in [s]$ comes with the weights δ_a, ρ_a, μ_a (or δ_a, d_a, μ_a). Equivalently, we can think that the information at each node is coded by the weights δ_a, ρ_a plus an additional label *global* or *local* according to where B_a is global or local, which prescribes the relation among δ_a, ρ_a and μ_a (or δ_a, d_a and μ_a).

Theorem 3.10. *The multiplicity μ of $\mathcal{J}_c$ at the origin equals*

$$\mu = \left(\prod_{a\in G\setminus S(G)} \rho_a\right)\left(\prod_{b\in S(G)} \mu_b\right). \tag{3.8}$$

Proof. We will prove formula 3.8 by induction in the number s of blocks. If $s = 1$, the system is irreducible and $\{1\} \in S(G)$ so the formula holds. Consider $s > 1$ and assume that the result is true for systems with $s - 1$ blocks. Let B be as in (3.7), set $n' := \nu_{s-1}$, where ν_{s-1} is as in (3.6), and consider the ideal $\mathcal{J}'_c := \langle p_1, \ldots, p_{n'}\rangle$, in the polynomial ring in the first n' variables. Clearly, $p_1, \ldots, p_{n'}$ is in normal and triangular form. Let G' be the corresponding graph; it is obtained by erasing from G the vertex s and all edges ending at s. By inductive hypothesis, we have that the multiplicity μ' of $\mathcal{J}'_c$ at $0'$ equals

$$\mu' = \left(\prod_{a\in G'\setminus S(G')} \rho_a\right)\left(\prod_{b\in S(G')} \mu_b\right). \tag{3.9}$$

The matrix B has the form

$$B = \left(\begin{array}{c|c} B' & 0 \\ \hline C & B_s \end{array}\right). \tag{3.10}$$

If the rectangular matrix C is identically zero, then the last $n - n'$ polynomials depend only on the last $n - n'$ variables, and we have that

$$\mu = \mu' \cdot \mu_s,$$

as wanted, since in this case $S(G) = S(G') \cup \{s\}$.

On the other hand, if C is not zero, it is possible to find a positive weight vector w such that the initial monomial $in_{-w}(p_j) = x_j^{r_j}$, for all $n' < j \leq n$. Consider any local order $\prec$ in $k[x_1, \ldots, x_n]$ refining the weight $-w$. Let $\{q_1, \ldots, q_t\}$ be a standard basis for the ideal $\mathcal{J}'_c$ with respect to the local order induced by $\prec$ in $k[x_1, \ldots, x_{n'}]$. Then, $\{q_1, \ldots, q_t, p_{n'+1}, \ldots, p_n\}$ is a standard basis for $\mathcal{J}_c$ relative to $\prec$ since, for every $i = 1, \ldots, t$, the leading monomials of the polynomial q_i is coprime with those of the p_j, $n' < j \leq n$, and, therefore, the weak normal form of the corresponding S-polynomial is 0 [13]. The corresponding initial ideal $L_\prec(\mathcal{J}_c)$ will be generated by some monomials in the first n' variables (which generate the initial ideal $L_{\prec'}(\mathcal{J}'_c)$) and the pure powers $x_j^{r_j}$ for all $j > n'$. Therefore, the multiplicity μ of $\mathcal{J}_c$ at 0 equals:

$$\begin{aligned} \dim_{\bar{k}}\left(\bar{k}[x_1 \ldots, x_n]/\mathcal{J}_c\right)_0 = \dim_{\bar{k}}\left(\bar{k}[x_1 \ldots, x_n]/L_\prec(\mathcal{J}_c)\right)_0 = \\ \dim_{\bar{k}}\left(\bar{k}[x_1 \ldots, x_{n'}]/L_{\prec'}(\mathcal{J}'_c)\right)_0 \cdot \dim_{\bar{k}}\left(\bar{k}[x_{n'+1} \ldots, x_n]/\langle x_{n'+1}^{r_{n'+1}} \ldots x_n^{r_n}\rangle\right)_0 = \\ \dim_{\bar{k}}\left(\bar{k}[x_1 \ldots, x_{n'}]/\mathcal{J}'_c\right)_0 \cdot \rho_k. \end{aligned}$$

In this case $s \notin S(G)$, and so $S(G) = S(G')$. Since the dimension of the local quotient by $\mathcal{J}'_c$ at the origin equals (3.9), we get that

$$\mu = \mu' \cdot \rho_s = \left(\prod_{a \in \setminus S(G)} \rho_a \right) \left(\prod_{b \in S(G)} \mu_b \right),$$

as wanted. □

Remark 3.11. Using Theorem 3.5 we can translate (3.8) as

$$\mu \;=\; \left(\prod_{a \in G_1} d_a \right) \left(\prod_{b \in G_2} \mu_b \right), \tag{3.11}$$

where G_1 is the set of nodes of G corresponding to the global, non-sources of G and G_2 is its complement.

We will also need the following terminology.

Definition 3.12. A vertex b of (the directed acyclic graph) G is said to be a *descendant* (respectively, a *direct descendant*) of the vertex a if there is a directed path (respectively, a directed edge) from a to b. A (directed) subgraph H of G is said to be *full* if, for any of its vertices j, all the descendants and all the directed paths starting from j also belong to H. The collection of full subgraphs of G will be denoted by $\mathcal{F}(G)$.

The empty subgraph is full and even if G is connected, a full subgraph H may be disconnected. Note also that a full subgraph is completely determined by its sources.

The following result refines the description given in Remark 2.7 of subsets $L \subset [n]$ with $\mu_L \neq 0$.

Proposition 3.13. *Let $p_1, \dots, p_n$ be a binomial complete intersection in normal and triangular form and $L \subset [n]$. Then $\mu_L = 0$ unless there exists a full subgraph H of G such that*

$$\prod_{a \notin H} \delta_a \neq 0 \tag{3.12}$$

and L coincides with the union of all the indices belonging to blocks which are vertices of H.

Proof. With the above notations, let $\lambda = (\lambda_1, \dots, \lambda_n) \in \mathbb{V}(\mathcal{J}_c)$ and $L = L(\lambda) = \{i \in [n] : \lambda_i = 0\}$. Set $H = \{a \in G : I_a \cap L \neq \emptyset\}$. If $a \in H$ then we may argue as in Lemma 3.4 to conclude that $I_a \subset L$. Suppose now that $a \in H$ and that (a, b) is an edge in G. Since $C_{ba} \neq 0$, there exists $i \in I_a$ and $j \in I_b$ such that $i \in \operatorname{supp}(\beta_j)$ and, consequently, $\lambda_j = 0$, i.e. $j \in I_b \cap L$, and $b \in H$. This shows that H is a full subgraph of H. The need for condition (3.12) was already noted in Remark 2.7. □

With notation as in Prop. 3.13, given a full subgraph $H \subset G$, we will denote by $L(H)$ the set of indices belonging to blocks associated with vertices of H.

Proposition 3.14. *Given a full subgraph H of G, the number $D_{L(H)}$ of points in $\mathbb{V}(\mathcal{J}_c) \cap \bar{k}^n_{L(H)}$ counted without multiplicity equals*

$$D_{L(H)} \;=\; \left(\prod_{a \notin H} \delta_a\right) \tag{3.13}$$

while the number $\mu_{L(H)}$ of points in $\mathbb{V}(\mathcal{J}_c) \cap \bar{k}^n_{L(H)}$ counted with multiplicity equals

$$\mu_{L(H)} \;=\; \left(\prod_{a \notin H} \delta_a\right) \left(\prod_{b \in H \setminus S(H)} \rho_b\right) \left(\prod_{e \in S(H)} \mu_e\right). \tag{3.14}$$

Proof. The first assertion follows easily from Proposition 3.13. In order to prove (3.14), let $\lambda \in \mathbb{V}(\mathcal{J}_c) \cap \bar{k}^n_{L(H)}$, write $\lambda = (\lambda^{(1)}, \dots, \lambda^{(s)})$ with $\lambda^{(a)} \in \bar{k}^{|I_a|}$ for all $a \in [s]$. Since H is a full subgraph, there are no edges starting at a node in H and ending at a node outside of H; i.e. $C_{ba} = 0$ for all $a \in H$ and $b \notin H$. Therefore, it is possible to relabel the variables and the binomials $p_1, \dots, p_n$ so that the system remains in normal form and satisfies that $a < b$ for all $a \notin H$ and $b \in H$. Thus, we may assume without loss of generality that $H = \{t+1, \dots, s\}$ and therefore $\lambda = (\lambda^{(1)}, \dots, \lambda^{(t)}, 0, \dots, 0)$ with $\lambda^{(a)} \in (\bar{k}^*)^{|I_a|}$ for $a = 1, \dots, t$. Equivalently,

$$\lambda = (\lambda', 0) \in (\bar{k}^*)^{n'} \times (\bar{k})^{n-n'} \;;\; n' := \nu_t\,.$$

We let x' stand for the first n' variables $x_1, \dots, x_{n'}$ and x'' for the remaining $n - n'$ variables. Then

$$\mathcal{J}'_c \;:=\; \langle p_1, \dots, p_{n'} \rangle \subset k[x']$$

and λ' is a simple zero of $\mathcal{J}'_c$. Hence $p_1, \dots, p_{n'}$ define the maximal ideal in the local ring $\left(\bar{k}[x']\right)_{\lambda'}$. We then have:

$$\begin{aligned}
\mu_\lambda \;&:=\; \dim_{\bar{k}} \left(\bar{k}[x]/\mathcal{J}_c\right)_\lambda \\
&=\; \dim_{\bar{k}} \left(\bar{k}[x]/\langle x_1 - \lambda_1, \dots, x_{n'} - \lambda_{n'}, p_{n'+1}, \dots, p_n \rangle\right)_\lambda \\
&=\; \dim_{\bar{k}} \left(\bar{k}[x'']/\langle p_{n'+1}(\lambda', x''), \dots, p_n(\lambda', x'') \rangle\right)_0 \\
&=\; \dim_{\bar{k}} \left(\bar{k}[x'']/\langle p_{N+1}(1, \dots, 1, x''), \dots, p_n(1, \dots, 1, x'') \rangle\right)_0 .
\end{aligned}$$

So, μ_λ equals the multiplicity at the origin $0 \in \bar{k}^{n-n'}$ of the system $\{\hat{p}_{n'+1}, \dots, \hat{p}_n\}$. Formula (3.14) now follows from Theorem 3.10, and the fact that the system $p_1, \dots, p_{n'}$ has $\delta_1 \cdots \delta_t$ simple solutions in $(\bar{k}^*)^{n'}$. □

The following explicit formulas for d and D follow by adding (3.13) and (3.14) over all full subgraphs of G.

Theorem 3.15. *Suppose that $p_1, \dots, p_n$ are in normal, triangular form. For generic parameters $c \in (k^*)^n$, the total number of solutions of the system $p_1(c; x) = \cdots = p_n(c; x) = 0$, counted without multiplicity, equals*

$$D \;=\; \sum_{H \in \mathcal{F}(G)} \left(\prod_{a \notin H} \delta_a\right), \tag{3.15}$$

and the total number of solutions counted with multiplicity equals

$$d \;=\; \sum_{H \in \mathcal{F}(G)} \left(\prod_{a \notin H} \delta_a\right) \left(\prod_{b \in H \setminus S(H)} \rho_b\right) \left(\prod_{e \in S(H)} \mu_e\right). \tag{3.16}$$

We end this section with a recursive formula to compute d. In order to state the following proposition we define, for $1 \leq r \leq s$, the binomial system $q^{(r)}$:

$$p_i(1, \dots, 1, x_{\nu_{r-1}+1}, \dots, x_n)\ , \quad i \in I_r \cup \cdots \cup I_s .$$

Note that the matrix associated with $q^{(r)}$ is:

$$B^{(r)} = \begin{pmatrix} B_r & 0 & \dots & 0 \\ C_{(r+1)r} & B_{r+1} & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ C_{sr} & C_{s(r+1)} & \dots & B_s \end{pmatrix} \tag{3.17}$$

Clearly if $p_1, \dots, p_n$ is in normal, triangular form, so is $q^{(r)}$. We denote by F_r the number of solutions in $\bar{k}^{n-\nu_{r-1}}$, counted with multiplicity, of the system q^r.

Proposition 3.16. *F_r is a function of $\{\delta_a, \mu_a, \rho_a\,;\ a = r, \dots, s\}$. The function F_r may be computed recursively as:*

$$F_s = d_s = \delta_s + \mu_s$$

$$F_r = \delta_r \cdot F_{r+1} + \mu_r \cdot F_{r+1}|_{\delta_b=0, \mu_b=\rho_b}, \tag{3.18}$$

where b runs over all indices in $\{r+1, \dots, s\}$ such that $C_{br} \neq 0$.

Proof. We may assume without loss of generality that $r = 1 < s$. Let G be the graph of B and $G^{(2)}$ the subgraph of G associated to the submatrix $B^{(2)}$ defined by (3.17).

Any full subgraph $H \in \mathcal{F}(G^{(2)})$ may be thought of as a full subgraph in G. We denote by $\mathcal{F}' \subset \mathcal{F}(G)$ the collection of such subgraphs. Clearly $\mathcal{F}'$ consists of all full subgraphs of G not containing the vertex 1. Let $\mathcal{F}''$ denote the complement of $\mathcal{F}'$ in $\mathcal{F}(G)$. Removing the vertex 1 from a subgraph $H \in \mathcal{F}''$ defines a full subgraph $H^{(2)}$ of $G^{(2)}$ with the property that no direct descendant of 1 in G may be in $G^{(2)} \setminus H^{(2)}$. Let us denote by $\mathcal{F}''(G^{(2)})$ the collection of such full subgraphs of $G^{(2)}$. We can write

$$F_1 \ = \sum_{H \in \mathcal{F}'} \mu_{L(H)} + \sum_{H \in \mathcal{F}''} \mu_{L(H)}. \tag{3.19}$$

Since, for $H \in \mathcal{F}'$, $1 \notin H$, in view of (3.14), the first sum may be computed as:

$$\sum_{H \in \mathcal{F}'} \mu_{L(H)} \ = \ \delta_1 \sum_{H \in \mathcal{F}(G^{(2)})} \mu_{L(H)} \ = \ \delta_1 F_2, \tag{3.20}$$

since $S(H)$ is the same whether we view H as a subgraph of G or of $G^{(2)}$.

Thus, in order to complete the proof we need to show that the second sum in (3.19) equals

$$\mu_1 \cdot F_2|_{\delta_b=0, \mu_b=\rho_b},$$

where b runs over all vertices in $G^{(2)}$ which are direct descendants of 1 in G. We note first of all, that setting $\delta_b = 0$ for all direct descendants b of 1 has the effect of restricting the sum in (3.16) to $\mathcal{F}''(G^{(2)})$. Moreover, given $H \in \mathcal{F}''$, let $H^{(2)}$ denote the full subgraph of $G^{(2)}$ obtained by removing the vertex 1 from H. Then $S(H^{(2)})$ consists of $S(H) \cap G^{(2)}$ together with all the direct descendants of 1 in H. This change may be accomplished by replacing μ_b by ρ_b whenever $b \in H^{(2)}$ is a direct descendant of 1 in H. Since $1 \in S(H)$ for all $H \in \mathcal{F}''$, we obtain the desired equality. $\square$

Example 3.17. We return to Example 2.13. We recall that the reduced system $\tilde{p}_1, \ldots, \tilde{p}_6$ is:

$$\tilde{p}_1 = u_1^5 - u_1^3 u_2; \quad \tilde{p}_2 = u_2^7 - u_1^8 u_2;$$
$$\tilde{p}_3 = u_3^2 - u_1^6 u_3^3; \quad \tilde{p}_4 = u_4^3 - u_1^2 u_2 u_5^2 u_6;$$
$$\tilde{p}_5 = u_5^2 - u_4^2 u_5; \quad \tilde{p}_6 = u_6^3 - u_1^3 u_4 u_5 u_6.$$

and, therefore, its associated matrix is

$$B = \begin{pmatrix} 2 & -1 & 0 & 0 & 0 & 0 \\ -8 & 6 & 0 & 0 & 0 & 0 \\ -6 & 0 & -1 & 0 & 0 & 0 \\ -2 & -1 & 0 & 3 & -2 & -1 \\ 0 & 0 & 0 & -2 & 1 & 0 \\ -3 & 0 & 0 & -1 & -1 & 2 \end{pmatrix}.$$

Therefore, the system is in normal, triangular form with blocks relative to the index sets $I_1 = \{1, 2\}$, $I_2 = \{3\}$, and $I_3 = \{4, 5, 6\}$. The block B_1 is global, while B_2 and B_3 are local. The graph G has 3 vertices $\{1, 2, 3\}$ and arrows from 1 to 2 and 1 to 3. Hence $S(G) = \{1\}$. The weights are:

$$\delta_1 = 4, \delta_2 = 1, \delta_3 = 5, \rho_1 = 35, \rho_2 = 2, \rho_3 = 18\,,$$

and, taking into account the local/global label, we get $\mu_1 = 31$, $\mu_2 = 2$, $\mu_3 = 18$.

We may now apply (3.8) to compute the multiplicity $\tilde{\mu}$ of $\langle \tilde{p}_1, \ldots, \tilde{p}_6 \rangle$ at the origin:

$$\tilde{\mu} = \mu_1 \cdot \rho_2 \cdot \rho_3 = 1116\,.$$

In order to compute $\tilde{d}$ we use the inductive procedure of Proposition 3.16. Since the subgraph with vertices $\{2, 3\}$ is disconnected we have:

$$F_2 = (\delta_2 + \mu_2) \cdot (\delta_3 + \mu_3).$$

Hence, $F_1 = \delta_1 \cdot (\delta_2 + \mu_2) \cdot (\delta_3 + \mu_3) + \mu_1 \cdot \rho_2 \cdot \rho_3$. This gives $\tilde{d} = 1392$. We note that this is far from the Bézout bound of 43740.

Using Lemma 3.1 and Theorem 3.2 we see that the total number of solutions for the original system $p_1, \ldots, p_8$ are given by $d = 2\tilde{d}$ and $\mu = 2\tilde{\mu}$. This values may be easily verified using a computer algebra system such as Singular [14].

Finally, we note that G has five full subgraphs with vertex sets: $\{1, 2, 3\}$, $\{2, 3\}$, $\{2\}$, $\{3\}$, and $\emptyset$. This means that there are five index sets $\tilde{L} \subset [6]$, such that $\mu_L \neq 0$. They are $\tilde{L}_1 = [6]$, $\tilde{L}_2 = \{3, 4, 5, 6\}$, $\tilde{L}_3 = \{3\}$, $\tilde{L}_4 = \{4, 5, 6\}$ and $\tilde{L}_5 = \emptyset$. The corresponding multiplicities are according to (3.14):

$$\mu_{\tilde{L}_1} = \tilde{\mu} = 1116,\ \mu_{\tilde{L}_2} = 144,\ \mu_{\tilde{L}_3} = 40,\ \mu_{\tilde{L}_4} = 72,\ \mu_{\tilde{L}_5} = \tilde{\delta} = 20.$$

Moreover, the total number of solutions counted without multiplicity is given by:

$$\tilde{D} = \delta_1 + \delta_1 \cdot \delta_3 + \delta_1 \cdot \delta_2 + \delta_1 \cdot \delta_2 \cdot \delta_3 = 48$$

This information may be lifted to the original system using the bijection $L \to \tilde{L}$ discussed before Theorem 3.2. We get that $\mu_L = 0$ except for the following subsets

$$L_1 = [8],\ L_2 = \{4, 5, 6, 7, 8\},\ L_3 = \{4\},\ L_4 = \{5, 6, 7, 8\},\ L_5 = \emptyset.$$

Once again, $\mu_{L_i} = 2\,\mu_{\tilde{L}_i}$.

4. Counting complexity

In this section we will study the counting complexity, in the sense of [31], of computing the numerical invariants d, D, δ, μ, and μ_L associated with a gci $p_1, \dots, p_n$.

We have already observed that we may decide in polynomial time whether $p_1, \dots, p_n$ is a gci and that the property of being a complete intersection is independent of the coefficients if $\det B \neq 0$. Moreover, if $p_1, \dots, p_n$ is a gci we may also transform it into normal and triangular form in quadratic time. Also, since a system with generic exponents is irreducible and satisfies $\det B \neq 0$, we may compute its invariants in time polynomial in n for any choice of coefficients. In the general case, we may compute δ, μ, and μ_L, for a particular choice of L, directly from the invariants δ_a, ρ_a, and μ_a associated with the diagonal blocks of the system. Thus, δ, μ, and μ_L may be computed in polynomial time as well. However, we will show below that the computation of d or D is a $\#P$-hard problem, and therefore it is at least as hard as an NP-hard problem [31]. In order to do this we begin by reversing the relationship between binomial systems and weighted acyclic directed graphs. We recall that to a binomial system $p_1, \dots, p_n$ in normal and triangular form we associate an acyclic directed graph G whose vertices $\{1, \dots, s\}$ correspond to the diagonal blocks of the associated matrix B and that each vertex has weights δ_a, ρ_a, $a \in [s]$, plus a label "local" or "global". In the first case we set $\mu_a = \rho_a$, while in the global case we set $\mu_a = \rho_a - \delta_a$. In any case $d_a = \delta_a + \mu_a$. The proof of the following proposition is straightforward.

Proposition 4.1. *Let $G = (V, E)$, $V = [s]$, be an acyclic directed graph, with weights $\delta_a, \rho_a \in \mathbb{Z}_{>0}$ and labels local/global attached to each vertex. Let μ_a and d_a be defined as above. Then, the system of binomials defined by*

$$p_a(x_1, \dots, x_s) = x_a^{d_a} - c_a \left(\prod_{(b,a) \in E} x_b \right) x_a^{\mu_a},$$

for all global vertices a, and

$$p_a(x_1, \dots, x_s) = x_a^{\mu_a} - c_a \left(\prod_{(b,a) \in E} x_b \right) x_a^{d_a},$$

for all local vertices a, has as weighted graph $(G, \delta_a, \rho_a, \mu_a)$.

Remark 4.2. The total number of solutions d and D of the system in Proposition 4.1 are given by (3.16) and (3.15). Note also that if a is a source of G, then we get $p_a = x_a^{d_a} - c_a x_a^{\mu_a}$ in the global case, and $p_a = x_a^{\mu_a} - c_a x_a^{d_a}$ in the local case. This is compatible with Remark 3.6.

In the particular case when all vertices $\{1, \dots, s\}$ of a directed acyclic graph G are local, and their weights are $\delta_a = 1$, $\rho_a = 1$, for all $a \in [s]$, the binomial system defined in Proposition 4.1 takes a very simple form:

$$p_a(x_1, \dots, x_s) = x_a - c_a \left(\prod_{(b,a) \in E} x_b \right) x_a^2, \; a = 1, \dots, s. \tag{4.1}$$

We will refer to this system as the *standard binomial system* associated with G.

Theorem 4.3. *Computing d and D for binomial complete intersections $p_1, \dots, p_n$ in normal, triangular form are $\#P$-hard problems.*

Proof. We will show that computing these invariants gives, for special binomial systems, the number of independent subsets of a bipartite graph G. Since, by [24], this is known to be a $\#P$-complete problem and any such problem is, by definition, $\#P$-hard the result will follow.

Let G be a bipartite graph with vertices $\{1, \ldots, s\}$. Let $p_1, \ldots, p_s$ be the standard binomial system of G as in (4.1). Then, for each full subgraph $H \subset G$ we have, by (3.14), that $\mu_{L(H)} = 1$. Hence, according to (3.16) and (3.15), both d and D are equal to the number of full subgraphs of G. But, as has been noted earlier, a full subgraph is completely determined by its sources and, for a bipartite graph G, a subset of vertices is the set of sources of a full subgraph H if and only if it is an independent subset of G. Thus, d and D agree with the number of independent subsets of G. □

Recall that a directed acyclic graph $G = (V, E)$ is called *transitive* if there is an edge $(a, b) \in E$ each time that there is a directed path from a to b, Transitive directed acyclic graphs are in correspondence with partial orders $\prec$ on V, where $a \prec b$ if and only if $(a, b) \in E$. Given a partial order $\prec$ on V, a subset A of V is called an antichain if given $a_1, a_2 \in A$, neither $a_1 \prec a_2$, nor $a_2 \prec a_1$. It is shown in [24] that counting the number of antichains in posets is a $\#P$-complete problem and, hence, $\#P$-hard. Given any directed acyclic graph $G = (V, E)$, it is possible to compute its transitive closure $G^+ = (V, E^+)$, in time $O(|V|^3)$ by the well known Floyd–Warshall's algorithm. It follows from (3.16) and (3.15) that d and D are the same for the standard binomial systems associated with G and with G^+.

Proposition 4.4. *The number of (simple) solutions, at a generic value of the coefficients, of the standard system associated with a directed acyclic graph G equals the number of antichains in the associated partial order.*

Proof. As in the proof of Theorem 4.3, for the standard binomial system of G we have $d = D$ and this number agrees with the number of full subgraphs of G. These subgraphs are determined by their sources, which correspond exactly to the antichains in the associated partial order on V. □

Although, as the previous results show, the problem of computing the total number of solutions for a general binomial system in normal and triangular form is $\#P$-hard, there are classes of binomial systems whose invariants may be computed in polynomial time. For example, if the graph is totally disconnected then $d = d_1 \cdots d_s = \prod_{i=1}^{s} (\delta_i + \mu_i)$. At the other extreme if G is a (complete) directed graph with vertices $\{1, \ldots, s\}$ and (b, a) is an edge of G for all $a, b \in [s]$ with $a < b$, then it is easy to see that there are only $s + 1$ full subgraphs of G and, consequently, the sums in (3.15) and (3.16), consist of $s + 1$ terms.

Even if the number of full subgraphs is exponential in s and G has few connected components, a bound on the number of local blocks guarantees that d can be computed in polynomial time in n. For instance, if all blocks are global, then B is an M-matrix and $p_1, \ldots, p_n$ is a Gröbner basis for a positive weight order, and so $d = \rho_1 \cdots \rho_s$. We end with the following "positive" complexity result.

Proposition 4.5. *Let $N \in \mathbb{Z}_{\geq 0}$. Assume $p_1, \ldots, p_n$ is in normal and triangular form with s blocks of which at most N are local. Then, there is a formula to compute the total multiplicity d with at most 2^N summands, each involving s products.*

Proof. Recall the notation in Proposition 3.16. We may write the formula F_r for the computation of the total number of solutions of the system $q^{(r)}$ purely in terms of δ_a and ρ_a by keeping track of the local/global character of each vertex and replacing μ_a by ρ_a if a is local and by $\rho_a - \delta_a$ in the case of a global vertex. Then, for a global vertex r, the recursion (3.18) becomes

$$F_r = \delta_r \cdot F_{r+1} + (\rho_r - \delta_r) \cdot F_{r+1}|_{\delta_a = 0}, \tag{4.2}$$

where a runs over all direct descendants of r. Let us write $F_{r+1} = F'_{r+1} + F''_{r+1}$, where F'_{r+1} consists of all summands containing a factor δ_a with a a direct descendant of 1. Hence, F'_{r+1} vanishes when we set such $\delta_a = 0$ and (4.2) becomes:

$$F_r = \delta_r \cdot (F'_{r+1} + F''_{r+1}) + (\rho_r - \delta_r) \cdot F''_{r+1} = \delta_r \cdot F'_{r+1} + \rho_r \cdot F''_{r+1}$$

and, consequently, the total number of summands does not change when adding a global vertex.

On the other hand, if B_r is local then (3.18) becomes

$$F_r = \delta_r \cdot F_{r+1} + \rho_r \cdot F_{r+1}|_{\delta_a = 0}$$

and the number of summands is, at worst, doubled. □

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