# Card shuffling, quantum mechanics and representation theory

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When G is infinite there it is believed that the choice of generators is not very important, and that "asymptotic" properties of the random walk depend only on the group. This has been proved in some cases, and has also supplied a number of exciting open problems.

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When G is the finite group  $S_n$ , the choice of generators matters dramatically. We are still very far from appreciating the richness that hides in the choice of generators.

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- The quantum Heisenberg ferromagnet
- A model for high rank linear groups

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There exists a universal constant C such that the diameter of any Cayley graph of  $S_n$  is smaller than  $n^C$ .

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- Helfgott-Seress-Żuk: For random generators the diameter is  $\leq Cn^2$ . (the conjectured diameter in this case is  $\approx n \log n$ )

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$$\lim_{n \to \infty} d(P_{(\frac{1}{2} - \epsilon)n \log n}, P_{\infty}) = 1.$$

where  $d(\cdot, \cdot)$  is the total variation distance and where  $P_{\infty}$  is the uniform measure on  $S_n$ .

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Generalizations for other conjugacy classes include Roichman (1996), Larsen-Shalev (2008) and Berestycki-Schramm-Zeitouni (2011).

#### Representation theory

For any function f on  $S_n$  we can define its "Fourier transform"  $\widehat{f}$ and it still satisfies that  $\widehat{f * g} = \widehat{f} \cdot \widehat{g}$  so the random walk probabilities satisfy  $\widehat{P}_t = (\widehat{P})^t$ .

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This is behind the analysis of Diaconis and Shahshahani. The actual values of  $\hat{f}$  go back to Frobenius (1901). More general results were obtained by Murnaghan (1937) and Nakayama (1940).

#### Take home message

# Representation theory is great if your generating set is a class function

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It has analogs of Parseval's formula and of the Fourier inversion formula.

Suppose  $\sigma_1 \geq \sigma_2 \geq \cdots \sigma_r$  and  $\sum \sigma_i = n$ . We call  $\sigma$  a partition of n and denote this by  $\sigma \vdash n$ . Partitions may be represented graphically as *Young diagrams*.

$$[5,1] = \boxed{\qquad} [3,2,1] = \boxed{\qquad} [2,1^3] = \boxed{\qquad}$$

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To relate the random walk on G to the random walk on the Cayley graph  $S_n$  we pass to continuous time. Define the (positive) laplacian  $\Delta(g) = |S|(\mathbf{1} - P)$  (here  $\Delta$ , P and  $\mathbf{1}$  are  $n! \times n!$  matrices). The probability to move from  $\sigma$  to  $\tau$  at time tis now given by  $e^{-t\Delta}(\sigma, \tau)$  (here this is matrix exponentiation).

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In continuous time X(i) is a continuous-time random walk on G for all i (dependent, of course). An equivalent point of view is that one puts marbles on the vertices of the graph and Poisson clocks on the edges. When a clock rings, exchange the marbles. In this formulation the process makes sense even on infinite graphs.

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- When the graph  $\mathbb{Z}^3$  is replaced by a tree there are results of Angel (2003) and Hammond (2013, preprint).

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The finite version was investigated with the cube replaced with the complete graph by Berestycki-Durrett (2006), Schramm (2005), Berestycki-K and Alon-K.

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#### Theorem (with Gil Alon, 2013)

Let G be any graph and let  $\lambda_1, \ldots, \lambda_{n-1}$  be the non-zero eigenvalues of the laplacian of continuous-time random walk on G. Let  $q_t$  be the probability that X(t) is a cycle of length n. Then

$$q_t = \frac{1}{n} \prod_{i=1}^{n-1} (1 - e^{-t\lambda_i})$$

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(as  $t \to \infty$ , the right-hand side converges to  $\frac{1}{n}$ , as it should. As  $t \to 0$  one gets a new proof of the matrix-tree theorem).

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Denote by Q the set of permutations which are on cycle of length n and write

$$q_t = \langle P_t, \mathbb{1}_Q \rangle = \langle \widehat{P}^t, \widehat{\mathbb{1}_Q} \rangle$$

Using Parseval's identity.

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Using Parseval's identity. Since  $\mathbb{1}_Q$  is a class function, its Fourier transform consists of scalar matrices. It is possible to calculate it explicitly and it is non-zero only on the hook-shaped diagrams  $[n-k, 1^k] =$ 

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Using Parseval's identity. Since  $\mathbb{1}_Q$  is a class function, its Fourier transform consists of scalar matrices. It is possible to calculate it explicitly and it is non-zero only on the hook-shaped diagrams  $[n-k, 1^k] = \square$ . The hook-shaped diagrams are very special: the dimension of the representation  $[n-k, 1^k]$  is  $\binom{n-1}{k}$  and  $\widehat{\Delta}([n-k, 1^k])$  may be diagonalized exactly and the eigenvalues are all sums of k-tuples of  $\lambda_i$  (Bacher, 1992).
#### Proof skeleton

#### Theorem

$$q_t = \frac{1}{n} \prod_{i=1}^{n-1} (1 - e^{-t\lambda_i})$$

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## Failure of the purely algebraic approach

Theorem

$$q_t = \frac{1}{n} \prod_{i=1}^{n-1} (1 - e^{-t\lambda_i})$$

Consider the event Q' that the permutation is a cycle of length n-1 and one fixed point. Then  $\widehat{\mathbb{1}_{Q'}}$  may still be calculated, but it is no longer supported on hook-shaped diagrams, but rather on diagrams with two rows and one column.

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Can we estimate  $\widehat{\Delta}(\rho)$  analytically? At least for the relevant  $\rho$  i.e. with two rows and one column?

Theorem (Caputo-Liggett-Richthammer, 2010)

If  $\rho \neq [n]$  then  $\lambda_1(\rho) \geq \lambda_1([n-1,1])$ .

 $(\lambda_1(\rho) \text{ stands for the smallest eigenvalue of } \widehat{\Delta}(\rho))$ . In other words, the smallest non-zero eigenvalue of  $\Delta$  is in the representation [n-1,1].

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#### Theorem (with Gil Alon, 2013)

If  $\rho$  has  $\leq \frac{1}{3}\sqrt{n}$  squares below the first row and  $\sigma$  has  $\leq \frac{1}{3}\sqrt{n}$  to the right of the leftmost column then  $\lambda_1(\rho) \leq \lambda_1(\sigma)$ .

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#### Take home message (speculative)

Representation theory is useful even when the generating set is not a class function, in combination with analytic methods.

# Thank you