

ACCURACY OF SEVERAL MULTIDIMENSIONAL REFINABLE DISTRIBUTIONS

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ABSTRACT. Compactly supported distributions f_1, \dots, f_r on \mathbf{R}^d are *refinable* if each f_i is a finite linear combination of the rescaled and translated distributions $f_j(Ax - k)$, where the translates k are taken along a lattice $\Gamma \subset \mathbf{R}^d$ and A is a *dilation matrix* that expansively maps Γ into itself. Refinable distributions satisfy a *refinement equation* $f(x) = \sum_{k \in \Lambda} c_k f(Ax - k)$, where Λ is a finite subset of Γ , the c_k are $r \times r$ matrices, and $f = (f_1, \dots, f_r)^T$. The *accuracy* of f is the highest degree p such that all multivariate polynomials q with $\deg(q) < p$ are exactly reproduced from linear combinations of translates of f_1, \dots, f_r along the lattice Γ . We determine the accuracy p from the matrices c_k . Moreover, we determine explicitly the coefficients $y_{\alpha,i}(k)$ such that $x^\alpha = \sum_{i=1}^r \sum_{k \in \Gamma} y_{\alpha,i}(k) f_i(x + k)$. These coefficients are multivariate polynomials $y_{\alpha,i}(x)$ of degree $|\alpha|$ evaluated at lattice points $k \in \Gamma$.

1. Introduction

A discrete set $\Gamma \subset \mathbf{R}^d$ is a *lattice* if it is the image of \mathbf{Z}^d under some nonsingular linear transformation. A $d \times d$ matrix A is *expansive* if every eigenvalue λ satisfies $|\lambda| > 1$. An expansive matrix A is a *dilation matrix* with respect to a lattice Γ if $A(\Gamma) \subset \Gamma$. Complex-valued functions f_1, \dots, f_r on \mathbf{R}^d are *refinable* with respect to A and Γ if each f_i equals a linear combination of the rescaled and translated functions $f_j(Ax - k)$. We shall only consider the case where the linear combinations are finite. Then the vector-valued function $f: \mathbf{R}^d \rightarrow \mathbf{C}^r$ defined by $f(x) = (f_1(x), \dots, f_r(x))^T$ satisfies a *refinement equation*, *dilation equation*, or *two-scale difference equation* of the form

$$f(x) = \sum_{k \in \Lambda} c_k f(Ax - k), \quad (1)$$

where Λ is a finite subset of Γ and the c_k are $r \times r$ matrices.

Refinable functions play important roles in several areas, including wavelet theory [Dau92] and subdivision schemes in approximation theory [CDM91]. A key goal is the determination of properties of a refinable $f: \mathbf{R}^d \rightarrow \mathbf{C}^r$ based on the coefficient mask $c = \{c_k\}_{k \in \Lambda}$. One fundamental property is the *accuracy* of f , the largest integer p such that every multivariate polynomial q with $\deg(q) < p$ lies in the shift-invariant space

$$S(f) = \left\{ \sum_{k \in \Gamma} \sum_{i=1}^r w_{k,i} f_i(x + k) : w_{k,i} \in \mathbf{C} \right\} = \left\{ \sum_{k \in \Gamma} w_k f(x + k) : w_k \in \mathbf{C}^{1 \times r} \right\} \quad (2)$$

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generated by f , where $\mathbf{C}^r = \mathbf{C}^{r \times 1}$ is the space of column vectors of length r and $\mathbf{C}^{1 \times r}$ is the space of row vectors of length r . We shall deal only with compactly supported f , in which case each series in (2) is well-defined for all choices of $w_{k,i}$.

The typical one-dimensional, single-function refinement equation has $d = 1$, $r = 1$, $\Gamma = \mathbf{Z}$, and $A = 2$. It is well-known that, with minor hypotheses, accuracy p holds in this case if and only if

$$\sum_{k=0}^N c_k = 2 \quad \text{and} \quad \sum_{k=0}^N (-1)^k k^j c_k = 0 \quad \text{for } j = 0, \dots, p-1.$$

The extension of these results to multiple functions, higher dimensions, general dilation matrices, and distributional solutions presents several difficulties. For the one-dimensional, multi-function case $d = 1$, $r \geq 1$, $\Gamma = \mathbf{Z}$, $A = 2$, Heil, Strang, and Strela [HSS96] and Plonka [Plo97] independently derived the ‘‘matrix sum rules’’ that characterize accuracy for integrable refinable functions. A previous paper [CHM98a] considered the higher-dimensional, multi-function case with an arbitrary dilation matrix A , again assuming the existence of an integrable, compactly supported solution to the refinement equation (1). Some similar results, for the case of diagonalizable A , were also derived in [Jng96].

The determination of whether a refinement equation has an integrable solution is considerably involved and complicated. However, every refinement equation such that the matrix $\Delta = \frac{1}{|\det(A)|} \sum_{k \in \Lambda} c_k$ has eigenvalues $\lambda_1 = \dots = \lambda_s = 1$, $|\lambda_{s+1}|, \dots, |\lambda_r| < 1$ with the eigenvalue 1 nondegenerate always has a compactly supported distributional solution (see Proposition 4.1).

In this paper we consider the general case of compactly supported distributional solutions to higher-dimensional, multi-function refinement equations with arbitrary dilation matrices. We present necessary and sufficient conditions for accuracy, and show that these conditions are equivalent to a finite set of finite linear equations. Our necessary condition requires only the additional hypothesis that translates of f be independent, while our sufficient condition does not require this hypothesis. For additional history and motivation of this problem, we refer to [CHM98a] and [JRZ97] and the references contained therein.

The basic notation for our paper is presented in Section 2. In Section 3 we present some results which apply to arbitrary distributions that are not necessarily refinable, and in Section 4 we present our main results on refinable distributions.

2. Notation

2.1. General Notation.

The Lebesgue measure of a set $E \subset \mathbf{R}^d$ is denoted $|E|$.

We use the standard multi-index notation $x^\alpha = x_1^{\alpha_1} \cdots x_d^{\alpha_d}$ for $\alpha = (\alpha_1, \dots, \alpha_d)$ and $x \in \mathbf{R}^d$. The degree of α is $|\alpha| = \alpha_1 + \dots + \alpha_d$. The number of multi-indices α of a given degree s is $d_s = \binom{s+d-1}{d-1}$. We write $\beta \leq \alpha$ if $\beta_i \leq \alpha_i$ for $i = 1, \dots, d$. If $\beta \leq \alpha$ then we set $\binom{\alpha}{\beta} = \binom{\alpha_1}{\beta_1} \cdots \binom{\alpha_d}{\beta_d}$, otherwise $\binom{\alpha}{\beta} = 0$.

A dilation matrix A necessarily has integer determinant. We set $m = |\det(A)|$. We let $d_1, \dots, d_m \in \Gamma$ be a *full set of digits*, i.e., a complete set of representatives of the order- m group $\Gamma/A(\Gamma)$. Then Γ is partitioned into the disjoint cosets

$$\Gamma_i = A(\Gamma) - d_i = \{Ak - d_i : k \in \Gamma\}.$$

Let $u_1, \dots, u_d \in \mathbf{R}^d$ be a set of *generators* for the lattice Γ , i.e.,

$$\Gamma = \{m_1 u_1 + \dots + m_d u_d : m_i \in \mathbf{Z}\}.$$

Then the rectangular parallelepiped

$$P = \{x_1 u_1 + \dots + x_d u_d : 0 \leq x_i < 1\} \quad (3)$$

is a *fundamental domain* for the group \mathbf{R}^d/Γ , and \mathbf{R}^d is partitioned into the disjoint sets $\{P+k\}_{k \in \Gamma}$.

We use the following generalized matrix notation. Let J and K be finite or countable index sets. Let $m_{j,k}$ be ordinary $r \times s$ matrices for $j \in J$ and $k \in K$. Then we say that $M = [m_{j,k}]_{j \in J, k \in K} \in (\mathbf{C}^{r \times s})^{J \times K}$ is a “ $J \times K$ matrix with $r \times s$ block entries.” If $N = [n_{k,\ell}]_{k \in K, \ell \in L} \in (\mathbf{C}^{s \times t})^{K \times L}$, then the product of the $J \times K$ matrix M with the $K \times L$ matrix N is the $J \times L$ matrix

$$MN = \left[\sum_{k \in K} m_{j,k} n_{k,\ell} \right]_{j \in J, \ell \in L}.$$

Most summations encountered in this paper will contain only finitely many nonzero terms.

A column vector is a $J \times 1$ matrix, which we denote by $v = [v_j]_{j \in J}$. The entries v_j may be scalars or $r \times s$ blocks. Analogously, a row vector is a $1 \times J$ matrix, which we denote by $u = (u_j)_{j \in J}$.

Integrals of the vector-valued function $f = (f_1, \dots, f_r)^\top$ are computed componentwise. If f is integrable then we define its Fourier transform by

$$\hat{f}(\omega) = \int_{\mathbf{R}^d} f(x) e^{-2\pi i x \cdot \omega} dx = \left(\int_{\mathbf{R}^d} f_1(x) e^{-2\pi i x \cdot \omega} dx, \dots, \int_{\mathbf{R}^d} f_r(x) e^{-2\pi i x \cdot \omega} dx \right)^\top.$$

2.2 Vector-Valued Distributions.

$C^\infty(\mathbf{R}^d)$ is the space of all infinitely differentiable functions on \mathbf{R}^d . The Schwartz class $\mathcal{S}(\mathbf{R}^d)$ contains all infinitely differentiable functions each of whose derivatives decay faster than the reciprocal of any polynomial. $C_c^\infty(\mathbf{R}^d)$ is the space of compactly supported infinitely differentiable functions. The topological dual of $C_c^\infty(\mathbf{R}^d)$ is the space of distributions $\mathcal{D}'(\mathbf{R}^d)$. The space of tempered distributions $\mathcal{S}'(\mathbf{R}^d)$ is the dual of $\mathcal{S}(\mathbf{R}^d)$. The space of compactly supported distributions $\mathcal{E}'(\mathbf{R}^d)$ is the dual of $C^\infty(\mathbf{R}^d)$. We write $\langle \varphi, g \rangle$ or $\langle \varphi(x), g(x) \rangle$ to denote the evaluation of a distribution g on a test function φ . The Fourier transform maps $\mathcal{S}(\mathbf{R}^d)$ into itself, and extends to $\mathcal{S}'(\mathbf{R}^d)$ by duality. The Paley–Wiener theorem for distributions implies that the Fourier transform of a compactly supported distribution is a continuous function on \mathbf{R}^d [Rud91, p. 198].

We let $C^\infty(\mathbf{R}^d, \mathbf{C}^r) = C^\infty(\mathbf{R}^d) \times \dots \times C^\infty(\mathbf{R}^d)$ denote the space of vector-valued functions $\varphi = (\varphi_1, \dots, \varphi_r)^\top$ with each component $\varphi_i \in C^\infty(\mathbf{R}^d)$. We use analogous notations for other cross products of spaces of test functions or distributions. If $\varphi = (\varphi_1, \dots, \varphi_r)^\top$ is a vector-valued test function and $f = (f_1, \dots, f_r)^\top$ is a vector-valued distribution, then we write

$$\langle \varphi, f \rangle = \begin{bmatrix} \langle \varphi_1, f_1 \rangle \\ \vdots \\ \langle \varphi_r, f_r \rangle \end{bmatrix}.$$

We allow a vector-valued distribution f to act on a scalar-valued test function φ by defining $\langle \varphi, f \rangle = (\langle \varphi, f_1 \rangle, \dots, \langle \varphi, f_r \rangle)^\top$ in this case.

Note that if $v = (v_1, \dots, v_r) \in \mathbf{C}^{1 \times r}$ is an ordinary row vector and $f = (f_1, \dots, f_r)^T \in \mathcal{E}'(\mathbf{R}^d, \mathbf{C}^r)$ is a compactly supported vector-valued distribution, then $vf = \sum_{i=1}^r v_i f_i \in \mathcal{E}'(\mathbf{R}^d)$ is the scalar-valued distribution defined by

$$\langle \varphi, vf \rangle = \sum_{i=1}^r \langle \varphi, v_i f_i \rangle, \quad \varphi \in C^\infty(\mathbf{R}^d).$$

In particular,

$$vf(0) = \sum_{i=1}^r v_i \hat{f}_i(0) = \sum_{i=1}^r v_i \langle 1, f_i \rangle = \left\langle 1, \sum_{i=1}^r v_i f_i \right\rangle = \langle 1, vf \rangle.$$

Given a compactly supported vector-valued distribution $f \in \mathcal{E}'(\mathbf{R}^d, \mathbf{C}^r)$ and given arbitrary row vectors $w_k \in \mathbf{C}^{1 \times r}$, the series $\sum_{k \in \Gamma} w_k f(x+k)$ defines a distribution in $\mathcal{D}'(\mathbf{R}^d)$ by the formula

$$\left\langle \varphi(x), \sum_{k \in \Gamma} w_k f(x+k) \right\rangle = \sum_{k \in \Gamma} w_k \langle \varphi(x-k), f(x) \rangle, \quad \varphi \in C_c^\infty(\mathbf{R}^d). \quad (4)$$

Since both f and φ are compactly supported, the right-hand side of (4) contains only finitely many nonzero terms. Letting $W = (w_k)_{k \in \Gamma}$ be the infinite row vector with block entries w_k , we write formally the infinite column vector of distributions

$$F(x) = [f(x+k)]_{k \in \Gamma},$$

and define WF to be the distribution in $\mathcal{D}'(\mathbf{R}^d)$ given by

$$WF(x) = \sum_{k \in \Gamma} w_k f(x+k).$$

Note that if the w_k grow in size at most polynomially, then $WF \in \mathcal{S}'(\mathbf{R}^d)$.

As a consequence of the preceding remarks, if $f \in \mathcal{E}'(\mathbf{R}^d, \mathbf{C}^r)$ then we can define the shift-invariant space $S(f)$ generated by f to be

$$S(f) = \left\{ \sum_{k \in \Gamma} w_k f(x+k) : w_k \in \mathbf{C}^{1 \times r} \right\} = \{WF : W \in (\mathbf{C}^{1 \times r})^{1 \times \Gamma}\}.$$

Note that $S(f)$ is a subspace of $\mathcal{D}'(\mathbf{R}^d)$. Since all polynomials are distributions, the definition of accuracy extends to distributions, i.e., we say that $f \in \mathcal{E}'(\mathbf{R}^d, \mathbf{C}^r)$ has accuracy p if each multivariate polynomial q with $\deg(q) < p$ lies in $S(f)$.

We say that translates of $f \in \mathcal{E}'(\mathbf{R}^d, \mathbf{C}^r)$ along Γ are *independent* if $WF(x) = \sum_{k \in \Gamma} w_k f(x+k)$ is the zero distribution if and only if $W = (w_k)_{k \in \Gamma}$ is the zero vector.

Finally, we say that a compactly supported vector-valued distribution $f \in \mathcal{E}'(\mathbf{R}^d, \mathbf{C}^r)$ is *refinable* if it is a solution of the refinement equation (1) in the sense of distributions, i.e., if

$$\langle \varphi, f \rangle = \sum_{k \in \Lambda} c_k \langle \varphi(x), f(Ax-k) \rangle = \frac{1}{m} \sum_{k \in \Lambda} c_k \langle \varphi(A^{-1}(x+k)), f(x) \rangle$$

for all $\varphi \in C^\infty(\mathbf{R}^d, \mathbf{C}^r)$. If we set $c_k = 0$ for $k \notin \Lambda$ and define L to be the $\Gamma \times \Gamma$ matrix

$$L = [c_{Ai-j}]_{i,j \in \Gamma},$$

then we can recast the refinement equation in the following form:

$$f(x) = \sum_{k \in \Lambda} c_k f(Ax - k) \quad \iff \quad F(x) = LF(Ax).$$

Additionally, if we define $B = (A^{-1})^T$ and let $M(\omega) = \frac{1}{m} \sum_{k \in \Lambda} c_k e^{-2\pi i k \cdot \omega}$ be the matrix-valued symbol of the refinement equation, then the refinement equation can be recast into the following form by applying the Fourier transform:

$$f(x) = \sum_{k \in \Lambda} c_k f(Ax - k) \quad \iff \quad \hat{f}(\omega) = M(B\omega) \hat{f}(B\omega).$$

In particular, if f is refinable and $\Delta = M(0) = \frac{1}{m} \sum_{k \in \Lambda} c_k$, then $\hat{f}(0) = M(0) \hat{f}(0) = \Delta \hat{f}(0)$. Therefore, $\hat{f}(0)$ is a right 1-eigenvector of Δ if $\hat{f}(0) \neq 0$.

2.3. Translation and Dilation of Multidimensional Polynomials.

We shall often deal with matrix-valued functions $u = [u_{j,k}]_{j \in J, k \in K}: \mathbf{R}^d \rightarrow \mathbf{C}^{J \times K}$ each of whose entries $u_{j,k}: \mathbf{R}^d \rightarrow \mathbf{C}$ is a polynomial. In this case, we refer to u as a *matrix of polynomials*. The degree of u is $\deg(u) = \max\{\deg(u_{j,k})\}_{j \in J, k \in K}$.

The number of monomials x^α of degree s is $d_s = \binom{s+d-1}{d-1}$. For a given degree $s \geq 0$, we collect the monomials of this degree together to form the vector of monomials $X_{[s]}: \mathbf{R}^d \rightarrow \mathbf{C}^{d_s}$ defined by

$$X_{[s]}(x) = [x^\alpha]_{|\alpha|=s}, \quad x \in \mathbf{R}^d.$$

For each integer $0 \leq t \leq s$, we define the matrix of polynomials $Q_{[s,t]}: \mathbf{R}^d \rightarrow \mathbf{C}^{d_s \times d_t}$ by

$$Q_{[s,t]}(y) = (-1)^{s-t} \left[\binom{\alpha}{\beta} y^{\alpha-\beta} \right]_{|\alpha|=s, |\beta|=t},$$

where $\binom{\alpha}{\beta} = \binom{\alpha_1}{\beta_1} \cdots \binom{\alpha_d}{\beta_d}$. Then, by [CHM98a], translation of $X_{[s]}(x)$ obeys the rule

$$X_{[s]}(x-y) = \sum_{t=0}^s Q_{[s,t]}(y) X_{[t]}(x).$$

Given any $d \times d$ matrix $Z = [z_{i,j}]_{i,j=1,\dots,d}$ and given $s \geq 0$, we let $Z_{[s]} = [z_{\alpha,\beta}^s]_{|\alpha|=s, |\beta|=s}$ be the $d_s \times d_s$ matrix whose scalar entries $z_{\alpha,\beta}^s$ are defined by the equation

$$\sum_{|\beta|=s} z_{\alpha,\beta}^s x^\beta = (Zx)^\alpha = \prod_{i=1}^d (z_{i,1}x_1 + \cdots + z_{i,d}x_d)^{\alpha_i}.$$

By [CHM98a], dilation of $X_{[s]}(x)$ by Z obeys the rule

$$X_{[s]}(Zx) = Z_{[s]} X_{[s]}(x).$$

The matrix $Z_{[s]}$ has a number of surprising properties. For example, if $\lambda = (\lambda_1, \dots, \lambda_d)^T$ is the vector consisting of all eigenvalues of Z , then $[\lambda^\alpha]_{|\alpha|=s}$ is the vector consisting of all eigenvalues of $Z_{[s]}$.

2.4. Some Special Matrices and Polynomial Functions.

Given a collection

$$\{v_\alpha = (v_{\alpha,1}, \dots, v_{\alpha,r}) \in \mathbf{C}^{1 \times r} : 0 \leq |\alpha| < p\}$$

of row vectors of length r , we shall associate a number of special matrices and functions.

We group the v_α by degree to form $d_s \times 1$ column vectors $v_{[s]} \in (\mathbf{C}^{1 \times r})^{d_s \times 1}$ with block entries that are the $1 \times r$ row vectors v_α , i.e.,

$$v_{[s]} = [v_\alpha]_{|\alpha|=s} = \begin{bmatrix} v_{\alpha_1,1} & \cdots & v_{\alpha_1,r} \\ \vdots & \ddots & \vdots \\ v_{\alpha_{d_s},1} & \cdots & v_{\alpha_{d_s},r} \end{bmatrix}.$$

Note that $v_{[0]} = [v_0] = v_0$.

For each α , we define a row vector of polynomials $y_\alpha: \mathbf{R}^d \rightarrow \mathbf{C}^{1 \times r}$ by

$$y_\alpha(x) = \sum_{0 \leq \beta \leq \alpha} (-1)^{|\alpha| - |\beta|} \binom{\alpha}{\beta} v_\beta x^{\alpha - \beta}.$$

If we write $y_\alpha(x) = (y_{\alpha,1}(x), \dots, y_{\alpha,r}(x))$, then the coefficients of the polynomial $y_{\alpha,i}$ are determined by the scalars $v_{\beta,i}$ for those β with $0 \leq \beta \leq \alpha$. Further, $\deg(y_\alpha) \leq |\alpha|$, and $\deg(y_\alpha) = |\alpha|$ if and only if $v_0 \neq 0$. In particular, y_0 is the constant polynomial $y_0(x) \equiv v_0$.

As with the vectors v_α , we collect the vectors of polynomials y_α by degree and arrange them as block entries in a column vector to form the matrix of polynomials $y_{[s]}: \mathbf{R}^d \rightarrow (\mathbf{C}^{1 \times r})^{d_s \times 1}$ defined by

$$y_{[s]}(x) = [y_\alpha(x)]_{|\alpha|=s} = \left[\sum_{t=0}^s \sum_{|\beta|=t} (-1)^{s-t} \binom{\alpha}{\beta} x^{\alpha-\beta} v_\beta \right]_{|\alpha|=s} = \sum_{t=0}^s Q_{[s,t]}(x) v_{[t]}.$$

We have $\deg(y_{[s]}) \leq s$, and $\deg(y_{[s]}) = s$ if and only if $v_0 \neq 0$.

Finally, for each x we collect the blocks $y_{[s]}(x+k)$ into an infinite row vector to form a function $Y_{[s]}: \mathbf{R}^d \rightarrow ((\mathbf{C}^{1 \times r})^{d_s \times 1})^{1 \times \Gamma}$. Specifically,

$$Y_{[s]}(x) = (y_{[s]}(x+k))_{k \in \Gamma}.$$

We adopt the convention that

$$Y_{[s]} = Y_{[s]}(0) = (y_{[s]}(k))_{k \in \Gamma}.$$

Thus $Y_{[s]}$ is the row vector of evaluations of the matrix of polynomials $y_{[s]}$ at lattice points.

The following fact on the behavior of the matrix of polynomials $y_{[s]}$ under translation will be useful.

Lemma 2.1. *Given a collection $\{v_\alpha \in \mathbf{C}^{1 \times r} : 0 \leq |\alpha| < p\}$ of row vectors, define the matrix of polynomials $y_{[s]}(x) = \sum_{t=0}^s Q_{[s,t]}(x) v_{[t]}$ as above. Then*

$$y_{[s]}(x+y) = \sum_{t=0}^s Q_{[s,t]}(y) y_{[t]}(x) \quad \text{and} \quad Y_{[s]}(x+y) = \sum_{t=0}^s Q_{[s,t]}(y) Y_{[t]}(x).$$

3. Results for Arbitrary Distributions

The following result states that if an arbitrary compactly supported vector-valued distribution f with independent translates has accuracy p , then the coefficients used to reconstruct polynomials from translates of f are themselves polynomials evaluated at lattice points. A direct proof of this result can be constructed from the proof of [CHM98a, Theorem 3.1] with distributional calculations replacing functional calculations. Alternatively, this theorem can be viewed as a statement about order of approximation, in which case proofs can be constructed by using the techniques of the papers [deB90], [BR92], [BDR94a], [BDR94b].

Theorem 3.1. *Assume that translates of $f \in \mathcal{E}'(\mathbf{R}^d, \mathbf{C}^r)$ along Γ are independent. If f has accuracy p , then there exists a collection $\{v_\alpha \in \mathbf{C}^{1 \times r} : 0 \leq |\alpha| < p\}$ of row vectors such that*

(i) $v_0 \neq 0$, and

$$(ii) X_{[s]}(x) = \sum_{k \in \Gamma} y_{[s]}(k) f(x+k) = Y_{[s]} F(x) \text{ for } 0 \leq s < p,$$

where $Y_{[s]} = (y_{[s]}(k))_{k \in \Gamma}$ is the row vector of evaluations at lattice points of the matrix of polynomials $y_{[s]}(x) = \sum_{t=0}^s Q_{[s,t]}(x) v_{[t]}$.

In particular, if q is any polynomial with $\deg(q) < p$, then there exists a unique row vector of polynomials $u_q: \mathbf{R}^d \rightarrow \mathbf{C}^{1 \times r}$, with $\deg(u_q) = \deg(u)$, such that $q(x) = \sum_{k \in \Gamma} u_q(k) f(x+k)$.

The following result (which also can be viewed as a consequence of Appell sequences), states that, regardless of whether f has accuracy p or not, if any monomial x^α can be reproduced from lattice translates of f using coefficients that are themselves polynomials evaluated at lattice points, then for each $0 \leq \beta \leq \alpha$ the monomial x^β can also be reproduced from translates of f . Moreover, the coefficients used to obtain x^β are the evaluations at lattice points of a constant times the $(\alpha - \beta)$ derivative of the coefficients used to obtain x^α .

Theorem 3.2. *Assume that $f \in \mathcal{E}'(\mathbf{R}^d, \mathbf{C}^r)$, and let α be any multi-index. If $u: \mathbf{R}^d \rightarrow \mathbf{C}^{1 \times r}$ is a row vector of polynomials such that*

$$x^\alpha = \sum_{k \in \Gamma} u(k) f(x+k),$$

then for each $0 \leq \beta \leq \alpha$,

$$x^\beta = C_\beta \sum_{k \in \Gamma} (D^{\alpha-\beta} u)(k) f(x+k),$$

where

$$D^\gamma u = \left(\frac{\partial^{|\gamma|}}{\partial x^\gamma} u_1, \dots, \frac{\partial^{|\gamma|}}{\partial x^\gamma} u_r \right) \quad \text{and} \quad C_\gamma = (-1)^{|\alpha-\gamma|} \frac{\gamma!}{\alpha!} = (-1)^{|\alpha-\gamma|} \frac{\gamma_1!}{\alpha_1!} \dots \frac{\gamma_d!}{\alpha_d!}.$$

Proof. Note that since $u(k)$ has polynomial growth, the series $\sum_{k \in \Gamma} u(k) f(x+k)$ defines a tempered

distribution. Fix any $\varphi \in \mathcal{S}(\mathbf{R}^d)$. Then

$$\begin{aligned}
\langle \varphi(x), (x + \ell)^\alpha \rangle &= \langle \varphi(x - \ell), x^\alpha \rangle \\
&= \left\langle \varphi(x - \ell), \sum_{k \in \Gamma} u(k) f(x + k) \right\rangle \\
&= \left\langle \varphi(x), \sum_{k \in \Gamma} u(k) f(x + \ell + k) \right\rangle \\
&= \left\langle \varphi(x), \sum_{k \in \Gamma} u(k - \ell) f(x + k) \right\rangle.
\end{aligned} \tag{5}$$

For each $y \in \mathbf{R}^d$, define $g_y, h_y \in \mathcal{S}'(\mathbf{R}^d)$ by

$$g_y(x) = (x + y)^\alpha \quad \text{and} \quad h_y(x) = \sum_{k \in \Gamma} u(k - y) f(x + k).$$

Given $\varphi \in \mathcal{S}(\mathbf{R}^d)$, the quantities

$$\begin{aligned}
P(y) &= \langle \varphi, g_y \rangle = \langle \varphi(x), (x + y)^\alpha \rangle, \\
Q(y) &= \langle \varphi, h_y \rangle = \left\langle \varphi(x), \sum_{k \in \Gamma} u(k - y) f(x + k) \right\rangle,
\end{aligned}$$

are both polynomials in the unknown y . From (5), we have $P(\ell) = Q(\ell)$ for every lattice point $\ell \in \Gamma$, and therefore $P(y) = Q(y)$ for every $y \in \mathbf{R}^d$.

Let e_j be the multi-index of degree 1 with a 1 in the j th coordinate and 0's elsewhere. Then for every $y \in \mathbf{R}^d$,

$$\langle \varphi(x), \alpha_j (x + y)^{\alpha - e_j} \rangle = \frac{\partial P}{\partial y_j}(y) = \frac{\partial Q}{\partial y_j}(y) = \sum_{k \in \Gamma} (-1) (D^{e_j} u)(k - y) \langle \varphi(x), f(x + k) \rangle.$$

Evaluating at $y = 0$, we have

$$\langle \varphi(x), \alpha_j x^{\alpha - e_j} \rangle = \sum_{k \in \Gamma} (-1) (D^{e_j} u)(k) \langle \varphi(x), f(x + k) \rangle.$$

Hence, distributionally,

$$\alpha_j x^{\alpha - e_j} = \sum_{k \in \Gamma} (-1) (D^{e_j} u)(k) f(x + k).$$

The general result then follows by iteration. \square

4. Results for Refinable Distributions

For the remainder of this paper we shall concentrate on compactly supported distributions which satisfy the refinement equation (1). We assume throughout that A is a dilation matrix with respect to a lattice Γ in \mathbf{R}^d , and that the coefficient mask $c = \{c_k\}_{k \in \Lambda}$ is a finite collection of $r \times r$ matrices.

We first note a mild condition on the coefficient mask which ensures that a compactly supported distributional solution of the refinement equation exists.

Proposition 4.1. Define $M(\omega) = \frac{1}{m} \sum_{k \in \Lambda} c_k e^{-2\pi i k \cdot \omega}$. If the matrix $\Delta = M(0) = \frac{1}{m} \sum_{k \in \Lambda} c_k$ has eigenvalues $\lambda_1 = \dots = \lambda_s = 1$, $|\lambda_{s+1}|, \dots, |\lambda_r| < 1$, with the eigenvalue 1 nondegenerate, then the following statements hold.

- (a) The infinite matrix product $P(\omega) = \prod_{j=1}^{\infty} M(B^j \omega)$ converges uniformly on compact sets to a continuous function with at most polynomial growth at infinity, where $B = (A^{-1})^T$.
- (b) If v is any right 1-eigenvector for Δ , then $\hat{f}(\omega) = P(\omega)v$ defines a distribution $f \in \mathcal{E}'(\mathbf{R}^d, \mathbf{C}^r)$ that satisfies the refinement equation (1), and $\hat{f}(0) = v$.
- (c) If $f \in \mathcal{E}'(\mathbf{R}^d, \mathbf{C}^r)$ is a distributional solution to the refinement equation (1), then $\hat{f}(\omega) = P(\omega) \hat{f}(0)$. In particular, if f is nontrivial then $\hat{f}(0) \neq 0$ and $\hat{f}(0)$ is a right 1-eigenvector for Δ .

Proof. (a) It was shown in [CHM98a, Theorem A.3] that the infinite matrix product $P(\omega)$ converges uniformly on compact sets to a continuous function. Since A is expansive, the matrix B is contractive. Hence, there exists a vector norm $|\cdot|$ on \mathbf{R}^d such that the corresponding matrix norm of B satisfies $|B| < 1$. Let $\|\cdot\|$ be any norm on \mathbf{C}^r , and define $K = \sup_{|\omega| \leq 1} \|P(\omega)\|$ and $R = \sup_{\omega \in \mathbf{R}^d} \|M(\omega)\|$. Set $\theta = 1/|B|$. Then

$$\sup_{|\omega| \leq \theta^n} \|P(\omega)\| \leq \sup_{|\omega| \leq \theta^n} \|P(B^n \omega)\| \|M(B^n \omega)\| \cdots \|M(B \omega)\| \leq K R^n.$$

Hence $\|P(\omega)\| \leq KR(1 + |\omega|)^{\log_{\theta} R}$, so P has at most polynomial growth at infinity.

(b) Let v be any right 1-eigenvector for Δ . By part (a), $\hat{f}(\omega) = P(\omega)v$ then defines a tempered distribution $f \in \mathcal{S}'(\mathbf{R}^d, \mathbf{C}^r)$. This distribution clearly satisfies $\hat{f}(\omega) = M(B\omega) \hat{f}(B\omega)$, so f is refinable. It therefore remains only to show that f is compactly supported. For each n , define $\mu_n \in \mathcal{S}'(\mathbf{R}^d, \mathbf{C}^r)$ by $\hat{\mu}_n(\omega) = (\prod_{j=1}^n M(B^j \omega))v$. Since $M(B^j \omega) = \frac{1}{m} \sum_{k \in \Lambda} c_k e^{-2\pi i k \cdot B^j \omega} = \frac{1}{m} \sum_{k \in \Lambda} c_k e^{-2\pi i A^{-j} k \cdot \omega}$, the entries of $\hat{\mu}_n(\omega)$ are finite linear combinations of the exponentials $e^{-2\pi i \ell \cdot \omega}$ with ℓ restricted to the discrete set $\Lambda_n = \sum_{j=1}^n A^{-j}(\Lambda)$. Since A^{-1} is contractive, there exists a compact set $\Omega \subset \mathbf{R}^d$ such that $\Lambda_n \subset \Omega$ for each n . Thus the entries of μ_n are finite linear combinations of point masses δ_{ℓ} with $\ell \in \Omega$. Hence $\text{supp}(\mu_n) \subset \Omega$. However, $\hat{\mu}_n(\omega) \rightarrow P(\omega)v = \hat{f}(\omega)$ uniformly on compact sets, so $\mu_n \rightarrow f$ weakly. Hence $\text{supp}(f) \subset \Omega$ as well.

(c) Note that $\hat{f}(\omega) = (\prod_{j=1}^n M(B^j \omega)) \hat{f}(B^j \omega)$ for each n , and that $\hat{f}(B^j \omega) \rightarrow \hat{f}(0)$ since B is contractive and \hat{f} is a continuous function. \square

In order to prove a result giving necessary and sufficient conditions for a refinable distribution to have accuracy p , we require the following ergodic-type lemma.

Lemma 4.2. Let $\mu \in \mathcal{S}'(\mathbf{R}^d)$. If there exists a $\lambda \in \mathbf{C}$ such that

$$\mu(Ax) = \lambda \mu(x) \quad \text{and} \quad \mu(x - \ell) = \mu(x), \quad \ell \in \Gamma,$$

then μ is a constant. Moreover, if $\lambda \neq 1$ then $\mu = 0$.

Proof. For each integer $j \in \mathbf{Z}$ and each lattice point $k \in \Gamma$, define $\mu_{j,k}(x) = \mu(A^j x - k)$. Then, by hypothesis, $\mu_{j,k}(x) = \lambda^j \mu(x)$. Let $B = (A^{-1})^T$. With $m = |\det(A)|$, we then have distributionally that

$$\lambda^j \hat{\mu}(\omega) = \hat{\mu}_{j,k}(\omega) = m^{-j} e^{-2\pi i k \cdot B^j \omega} \hat{\mu}(B^j \omega).$$

Choose now any $\varphi \in \mathcal{S}(\mathbf{R}^d)$ such that $\hat{\varphi}$ is compactly supported. Then $\nu_j(\omega) = \hat{\varphi}((A^T)^j \omega) \hat{\mu}(\omega)$ is a compactly supported tempered distribution. Moreover, $\text{supp}(\nu_j) \subset \text{supp}(\hat{\varphi}((A^T)^j \omega))$. Since A is expansive, by choosing j large enough we will have $\text{supp}(\nu_j)$ contained in the rectangular parallelepiped $R = \{x_1 u_1 + \cdots + x_d u_d : -1/2 \leq x_i < 1/2\}$, which is a fundamental domain for Γ . In this case, ν_j is completely determined by the values $\hat{\nu}_j(k)$ for $k \in \Gamma$. Since ν_j is compactly supported, we can compute these values as follows.

$$\begin{aligned}
\hat{\nu}_j(k) &= \langle e^{-2\pi i k \cdot \omega}, \nu_j(\omega) \rangle \\
&= \langle e^{-2\pi i k \cdot \omega}, \hat{\varphi}((A^T)^j \omega) \hat{\mu}(\omega) \rangle \\
&= \langle \hat{\varphi}((A^T)^j \omega), e^{-2\pi i k \cdot \omega} \hat{\mu}(\omega) \rangle \\
&= \langle \hat{\varphi}(\omega), m^{-j} e^{-2\pi i k \cdot (A^T)^{-j} \omega} \hat{\mu}((A^T)^{-j} \omega) \rangle \\
&= \langle \hat{\varphi}, \hat{\mu}_{j,k} \rangle \\
&= \langle \hat{\varphi}, \lambda^j \hat{\mu} \rangle = C.
\end{aligned}$$

Thus $\hat{\nu}_j(k)$ is a constant independent of k , so $\nu_j = C \delta$. Since $\nu_j(\omega) = \hat{\varphi}((A^T)^j \omega) \hat{\mu}(\omega)$, it follows that $\text{supp}(\hat{\mu}) = \{0\}$. As a consequence, μ must be a polynomial [Rud91, p. 194]. However, μ is Γ -periodic by hypothesis, so this implies that μ must be a constant. Finally, since $\mu(Ax) = \lambda \mu(x)$, this constant must be zero if $\lambda \neq 1$. \square

The following result gives necessary and/or sufficient conditions for a refinable distribution to have accuracy p . Recall that P is the fundamental domain for the lattice Γ defined in (3).

Theorem 4.3. *Assume that $f \in \mathcal{E}'(\mathbf{R}^d, \mathbf{C}^r)$ is a distributional solution of the refinement equation (1). Consider the following statements.*

(I) *f has accuracy p .*

(II) *There exists a collection of row vectors $\{v_\alpha \in \mathbf{C}^{1 \times r} : 0 \leq |\alpha| < p\}$ such that*

(i) *$v_0 \hat{f}(0) \neq 0$, and*

(ii) *$Y_{[s]} = A_{[s]} Y_{[s]} L$ for $0 \leq s < p$,*

where $Y_{[s]} = (y_{[s]}(k))_{k \in \Gamma}$ is the row vector of evaluations at lattice points of the matrix of polynomials $y_{[s]}(x) = \sum_{t=0}^s Q_{[s,t]}(x) v_{[t]}$.

Then we have the following.

(a) *If translates of f along Γ are independent, then statement (I) implies statement (II).*

(b) *Statement (II) implies statement (I). Moreover, in this case, after scaling the vectors v_α by the constant $C = (v_0 \hat{f}(0))^{-1} |P|$, we have*

$$X_{[s]}(x) = \sum_{k \in \Gamma} y_{[s]}(k) f(x+k) = Y_{[s]} F(x), \quad 0 \leq s < p.$$

Proof. (a) Assume that f has accuracy p and that translates of f along Γ are independent. Then Theorem 3.1 implies that there exist row vectors $\{v_\alpha \in \mathbf{C}^{1 \times r} : 0 \leq |\alpha| < p\}$ such that

$$X_{[s]}(x) = \sum_{k \in \Gamma} y_{[s]}(k) f(x+k) = Y_{[s]} F(x), \quad 0 \leq s < p. \quad (6)$$

Combining this with the refinement equation $F(x) = L F(Ax)$ and with the definition of $A_{[s]}$, we have

$$Y_{[s]} F(Ax) = X_{[s]}(Ax) = A_{[s]} X_{[s]}(x) = A_{[s]} Y_{[s]} F(x) = A_{[s]} Y_{[s]} L F(Ax).$$

Since translates of f are independent, this implies that $Y_{[s]}^T = A_{[s]} Y_{[s]}^T L$ for $0 \leq s < p$, which proves (ii).

To prove (i), consider in (6) the case $s = 0$. Since $y_{[0]}(k) = v_0$ for every k , we have

$$1 = x^0 = X_{[0]}(x) = \sum_{k \in \Gamma} v_0 f(x + k).$$

Recall that the rectangular parallelepiped $P = \{x_1 u_1 + \cdots + x_d u_d : 0 \leq x_i < 1\}$ is a fundamental domain for Γ . Let $\varphi \in C_c^\infty(\mathbf{R}^d)$ be a nonnegative, compactly supported function such that $\sum_{k \in \Gamma} \varphi(x - k) = 1$. Then, necessarily, $\int \varphi(x) dx = |P|$. Therefore,

$$\begin{aligned} v_0 \hat{f}(0) &= \langle 1, v_0 f \rangle \\ &= \left\langle \sum_{k \in \Gamma} \varphi(x - k), v_0 f(x) \right\rangle \\ &= \left\langle \varphi(x), \sum_{k \in \Gamma} v_0 f(x + k) \right\rangle \\ &= \langle \varphi, 1 \rangle \\ &= \int_{\mathbf{R}^d} \varphi(x) dx \\ &= |P| \neq 0, \end{aligned}$$

which completes the proof.

(b) Assume that statement (II) holds. For each $0 \leq s < p$, define the vector-valued distribution $G_{[s]} \in \mathcal{S}'(\mathbf{R}^d, \mathbf{C}^{d_s})$ by

$$G_{[s]}(x) = \sum_{k \in \Gamma} y_{[s]}(k) f(x + k) = Y_{[s]} F(x).$$

Using the equation $Y_{[s]} = A_{[s]} Y_{[s]} L$ and the refinement equation $L F(Ax) = F(x)$, we have

$$G_{[s]}(Ax) = Y_{[s]} F(Ax) = A_{[s]} Y_{[s]} L F(Ax) = A_{[s]} Y_{[s]} F(x) = A_{[s]} G_{[s]}(x). \quad (7)$$

We will show by induction that there is a constant C independent of s so that $G_{[s]}(x) = C X_{[s]}(x)$ for $0 \leq s < p$, and we will show that the explicit value of C is $C = (v_0 \hat{f}(0)) |P|^{-1}$.

Consider the case $s = 0$. We have $d_0 = 1$, so $G_{[0]} \in \mathcal{S}'(\mathbf{R}^d)$. Since $A_{[0]} = 1$, we have by (7) that $G_{[0]}(Ax) = G_{[0]}(x)$, and since $y_{[0]}(k) = v_0$ for every k , we also have that $G_{[0]}(x - \ell) = \sum_{k \in \Gamma} v_0 f(x - \ell + k) = G_{[0]}(x)$ for every $\ell \in \Gamma$. Lemma 4.2 therefore implies that $G_{[0]}$ is a constant C . To evaluate this constant explicitly, fix a nonnegative, compactly supported function $\varphi \in \mathcal{S}(\mathbf{R}^d)$ with $\text{supp}(\varphi) \subset P$ such that $\sum_{k \in \Gamma} \varphi(x - k) = 1$. Then $\int \varphi(x) dx = |P|$, so

$$\begin{aligned} C |P| &= \langle \varphi, C \rangle \\ &= \langle \varphi, G_{[0]} \rangle \\ &= \left\langle \varphi(x), \sum_{k \in \Gamma} v_0 f(x + k) \right\rangle \end{aligned}$$

$$\begin{aligned}
&= \left\langle \sum_{k \in \Gamma} \varphi(x - k), v_0 f(x) \right\rangle \\
&= \langle 1, v_0 f \rangle \\
&= v_0 \hat{f}(0) \neq 0.
\end{aligned}$$

Suppose now, inductively, that $G_{[t]}(x) = C X_{[t]}(x)$ for $0 \leq t < s$. Then, recalling the notation $Y_{[s]}(x) = (y_{[s]}(x + k))_{k \in \Gamma}$, we have

$$\begin{aligned}
G_{[s]}(x - \ell) &= Y_{[s]} F(x - \ell) \\
&= Y_{[s]}(\ell) F(x) \\
&= \sum_{t=0}^s Q_{[s,t]}(\ell) Y_{[t]} F(x) && \text{by Lemma 2.1} \\
&= \sum_{t=0}^s Q_{[s,t]}(\ell) G_{[t]}(x) && \text{by definition of } G_{[t]} \\
&= Q_{[s,s]}(\ell) G_{[s]}(x) + \sum_{t=0}^{s-1} Q_{[s,t]}(\ell) G_{[t]}(x) \\
&= Q_{[s,s]}(\ell) G_{[s]}(x) + C \sum_{t=0}^{s-1} Q_{[s,t]}(\ell) X_{[t]}(x) && \text{by induction} \\
&= Q_{[s,s]}(\ell) G_{[s]}(x) + C \sum_{t=0}^s Q_{[s,t]}(\ell) X_{[t]}(x) - C Q_{[s,s]}(\ell) X_{[s]}(x) \\
&= G_{[s]}(x) + C X_{[s]}(x - \ell) - C X_{[s]}(x) && \text{by definition of } Q_{[s,t]}.
\end{aligned}$$

Defining $H_{[s]}(x) = G_{[s]}(x) - C X_{[s]}(x)$, it therefore follows from the preceding calculation that

$$H_{[s]}(x - \ell) = H_{[s]}(x), \quad \ell \in \Gamma. \quad (8)$$

In addition, it follows from (7) that $H_{[s]}$ also satisfies

$$H_{[s]}(Ax) = A_{[s]} H_{[s]}(x).$$

We will now invoke Lemma 4.2 to show that each of the entries of $H_{[s]}$ is zero. As in [CHM98a, Lemma 4.2], by choosing an appropriate ordering of the monomials x^α of degree s , we may assume that the $d_s \times d_s$ matrix $A_{[s]}$ is lower-triangular. Let $\lambda_1, \dots, \lambda_{d_s}$ be the eigenvalues of $A_{[s]}$. Let $h_i \in \mathcal{S}'(\mathbf{R}^d)$ be the i th component of $H_{[s]}$, and let h_n be the first nonzero component, so that $H_{[s]} = (0, \dots, 0, h_n, \dots, h_{d_s})^\top$. By (8), we have $h_n(x - \ell) = h_n(x)$ for $\ell \in \Gamma$. Moreover, since $A_{[s]}$ is lower-triangular with $\lambda_1, \dots, \lambda_{d_s}$ on its diagonal and since $H_{[s]}(Ax) = A_{[s]} H_{[s]}(x)$, we also have $h_n(Ax) = \lambda_n h_n(x)$. Since A is expansive and since $s > 0$, it follows from [CHM98a, Lemma 4.2] that $A_{[s]}$ is also expansive. Therefore $|\lambda_n| > 1$, so Lemma 4.2 implies that $h_n = 0$. Hence $H_{[s]} = 0$, whence $G_{[s]}(x) = C X_{[s]}(x)$. This completes the proof. \square

Theorem 4.3 allows us to determine the accuracy of a compactly supported refinable distribution. The condition $Y_{[s]} = A_{[s]} Y_{[s]} L$ in statement (II) of Theorem 4.3 involves only the coefficients c_k and the structure of the matrix L . In particular, it does not depend on whether the solution of the refinement equation is a function or a distribution. It was shown in [CHM98a] that the seemingly infinite set of conditions on the vectors v_α given by the statement $Y_{[s]} = A_{[s]} Y_{[s]} L$ is in fact equivalent to a finite system of finite linear equations. This equivalence is quoted here for completeness, since it will be used for further results below.

Theorem 4.4. *Let $m = |\det(A)|$, and let $d_1, \dots, d_m \in \Gamma$ be a full set of digits. Set $\Gamma_i = A(\Gamma) - d_i$. Given a collection $\{v_\alpha \in \mathbf{C}^{1 \times r} : 0 \leq |\alpha| < p\}$ of row vectors, let $y_{[s]}(x) = \sum_{t=0}^s Q_{[s,t]}(x) v_{[t]}$ be the associated matrix of polynomials and let $Y_{[s]} = (y_{[s]}(k))_{k \in \Gamma}$ be the row vector of evaluations of these polynomials at lattice points.*

If $v_0 \neq 0$, then the following statements are equivalent.

- (a) $Y_{[p-1]} = A_{[p-1]} Y_{[p-1]} L$.
- (b) $Y_{[s]} = A_{[s]} Y_{[s]} L$ for $0 \leq s < p$.
- (c) $v_{[s]} = \sum_{k \in \Gamma_i} \sum_{t=0}^s Q_{[s,t]}(k) A_{[t]} v_{[t]} c_k$ for $0 \leq s < p$ and $i = 1, \dots, m$.

The test for accuracy in statement (II) of Theorem 4.3 includes the condition that $v_0 \hat{f}(0) \neq 0$. We can formulate the following sufficient conditions so that $v_0 \neq 0$ implies $v_0 \hat{f}(0) \neq 0$.

Theorem 4.5. *Assume that $f \in \mathcal{E}'(\mathbf{R}^d, \mathbf{C}^r)$ is a distributional solution of the refinement equation (1.1). Let $m = |\det(A)|$, and let $d_1, \dots, d_m \in \Gamma$ be a full set of digits. Assume that $v_0 \in \mathbf{C}^{1 \times r}$ satisfies statement (c) in Theorem 4.4 for the case $s = 0$, i.e.,*

$$v_0 = v_0 \sum_{k \in \Gamma_i} c_k, \quad i = 1, \dots, m.$$

If $v_0 \neq 0$, then either of the following two conditions is sufficient to imply that $v_0 \hat{f}(0) \neq 0$, and therefore that f has accuracy at least $p = 1$:

- (a) *translates of f along Γ are independent, or*
- (b) *the matrix $\Delta = \frac{1}{m} \sum_{k \in \Lambda} c_k$ satisfies “condition E(1),” i.e., it has eigenvalues $\lambda_1 = 1$ and $|\lambda_2|, \dots, |\lambda_r| < 1$.*

Proof. (a) Define $G_{[0]}(x) = v_0 \sum_{k \in \Gamma} f(x+k)$. Then the argument of the proof of Theorem 4.3(b) shows that $G_{[0]}(x) = C$ a.e., with $C = (v_0 \hat{f}(0)) |P|^{-1}$. Hence $v_0 \hat{f}(0) \neq 0$ if and only if $C \neq 0$. However, if translates of f along Γ are independent, then we must have $C \neq 0$ since $C = \sum_{k \in \Gamma} v_0 f(x+k)$ and $v_0 \neq 0$.

(b) Assume that $\Delta = \frac{1}{m} \sum_{k \in \Lambda} c_k$ has eigenvalues $\lambda_1 = 1$ and $|\lambda_2|, \dots, |\lambda_r| < 1$. Then, by Proposition 4.1(b), $\hat{f}(0)$ is the right 1-eigenvector for Δ . On the other hand, since $v_0 = v_0 \sum_{k \in \Gamma_i} c_k$ and since Γ is the disjoint union of the Γ_i , we have

$$v_0 = v_0 \frac{1}{m} \sum_{i=1}^m \sum_{k \in \Gamma_i} c_k = v_0 \Delta.$$

Hence v_0 is the left 1-eigenvector for Δ . Since the dot product of the left and right 1-eigenvectors must be nonzero when 1 is a simple eigenvalue, we have $v_0 \hat{f}(0) \neq 0$. \square

Thus, if either of these two conditions are satisfied, in order to determine the accuracy of f we can use any of the three equivalent statements of Theorem 4.4. Moreover, the following result implies that test (b) of Theorem 4.4 imposes a necessary condition on the left eigenvalues of L .

Proposition 4.6. *Let $\lambda = (\lambda_1, \dots, \lambda_d)^T$ be the vector of all eigenvalues of A . If there exist row vectors $Y_{[s]} \in ((\mathbf{C}^{1 \times r})^{d_s \times 1})^{1 \times \Gamma}$ such that $Y_{[s]} = A_{[s]} Y_{[s]} L$ for $0 \leq s < p$, then $\lambda^{-\alpha}$ is a left eigenvalue for L for each multi-index α with $0 \leq |\alpha| < p$.*

Proof. By [CHM98a, Lemma 4.2], the eigenvalues of $A_{[s]}$ are $[\lambda^\alpha]_{|\alpha|=s}$. Let S be such that $J = S^{-1} A_{[s]} S$ is in Jordan form. If J is diagonal, then its diagonal entries are λ^α . Therefore, by thinking of $Z_{[s]} = S^{-1} Y_{[s]} = [z_\alpha]_{|\alpha|=s}$ as having “rows” $z_\alpha = (z_\alpha(k))_{k \in \Gamma} \in (\mathbf{C}^{1 \times r})^{1 \times \Gamma}$, we can compute

$$[z_\alpha]_{|\alpha|=s} = Z_{[s]} = S^{-1} Y_{[s]} = S^{-1} A_{[s]} S S^{-1} Y_{[s]} L = J Z_{[s]} L = [\lambda^\alpha z_\alpha L]_{|\alpha|=s}.$$

Thus z_α is a left $\lambda^{-\alpha}$ -eigenvector for L for each $|\alpha| = s$. If J is not diagonal, then for each distinct value of λ^α there is still at least one z_α such that $z_\alpha = \lambda^\alpha z_\alpha L$, so each of the distinct values of $\lambda^{-\alpha}$ is still a left eigenvalue for L . \square

Considering Theorem 4.3, Theorem 4.4, and Proposition 4.6 together, we see that if f is to have accuracy p , then $\lambda^{-\alpha}$ must be a left eigenvalue for L for each $0 \leq |\alpha| < p$. An example from [JRZ97] shows that even in the case $d = 1$, $r = 1$, the existence of such eigenvalues alone is not sufficient to imply accuracy for f ; the corresponding left eigenvectors must have the polynomial structure specified in Theorem 4.3.

Since L is an infinite matrix, it is conceivable that the determination of its eigenvalues could be a difficult task. In fact, the eigenvalues and eigenvectors of L are completely determined by a particular finite submatrix of L . This was shown by Jia for the one-dimensional, single-function case in [Jia96]. The higher-dimensional, single-function case was considered in [Jia98], and the one-dimensional, multi-function case was discussed in [JRZ97]. Because the characterization of the eigenvalues of L leads to an alternative test for the accuracy of a refinable distribution, we briefly sketch the extension of these ideas to the general higher-dimensional, multi-function setting of this paper.

Define the support of a column vector $a = [a_k]_{k \in \Gamma}$ to be $\text{supp}(a) = \{k \in \Gamma : a_k \neq 0\}$, and set

$$\begin{aligned} \ell(\Gamma) &= \{a = [a_k]_{k \in \Gamma} : a_k \in \mathbf{C}^{r \times 1}\} = (\mathbf{C}^{r \times 1})^{\Gamma \times 1}, \\ \ell_c(\Gamma) &= \{a \in \ell(\Gamma) : \text{supp}(a) \text{ is finite}\}. \end{aligned}$$

For each nonempty $\Omega \subset \Gamma$ define

$$\ell(\Omega) = \{a \in \ell(\Gamma) : \text{supp}(a) \subset \Omega\},$$

and define $\ell(\emptyset) = \{0\}$. We will study a class of finite sets Ω for which $\ell(\Omega)$ is right-invariant under L . In particular, we will show in Lemma 4.7 that $\ell(\Omega)$ is invariant if

$$A^{-1}(\Omega + \Lambda) \cap \Gamma \subset \Omega. \tag{9}$$

We will say that a finite nonempty set $\Omega \subset \Gamma$ is *admissible* if (9) holds.

The language of Iterated Function Systems (IFS) is convenient for discussing the properties of admissible sets. Let $\mathcal{H}(\mathbf{R}^d)$ be the metric space of all nonempty compact subsets of \mathbf{R}^d under the Hausdorff metric. For each $k \in \Lambda$, define $w_k: \mathbf{R}^d \rightarrow \mathbf{R}^d$ by $w_k(x) = A^{-1}(x + k)$. Then define $w_\Lambda: \mathcal{H}(\mathbf{R}^d) \rightarrow \mathcal{H}(\mathbf{R}^d)$ by

$$w_\Lambda(K) = \bigcup_{k \in \Lambda} w_k(K) = A^{-1}(K + \Lambda), \quad K \in \mathcal{H}(\mathbf{R}^d).$$

Note that:

$$\Omega \subset \Gamma \text{ is admissible} \iff w_\Lambda(\Omega) \cap \Gamma \subset \Omega.$$

Since A is expansive, there exists a vector norm $\|\cdot\|$ on \mathbf{R}^d such that $\|A^{-1}\| < 1$. Therefore each w_k is a contractive mapping on \mathbf{R}^d , and as a consequence it can be shown that w_Λ is a contractive mapping on $\mathcal{H}(\mathbf{R}^d)$. By the Contraction Mapping Theorem, there must therefore exist a unique nonempty compact set $K_\Lambda \subset \mathbf{R}^d$ such that $w_\Lambda(K_\Lambda) = K_\Lambda$, i.e., such that $A^{-1}(K_\Lambda + \Lambda) = K_\Lambda$. In fact,

$$K_\Lambda = \sum_{n=1}^{\infty} A^{-n}(\Lambda) = \left\{ \sum_{n=1}^{\infty} A^{-n} \lambda_n : \lambda_n \in \Lambda \right\}. \quad (10)$$

The set K_Λ is called the *attractor* of the iterated function system generated by $\{w_k\}_{k \in \Lambda}$ [Hut81]. It can be shown that if f is compactly supported and satisfies the refinement equation (1), then $\text{supp}(f) \subset K_\Lambda$ [CHM98b].

We now prove some basic properties of admissible sets. In particular, we show that the set

$$\Omega_\Lambda = K_\Lambda \cap \Gamma$$

is admissible, and possesses some important special properties among the class of all admissible sets. For the remainder of this article, we will let $\|\cdot\|$ denote any vector norm on \mathbf{R}^d such that $\|A^{-1}\| < 1$, and we let

$$B(\rho) = \{x \in \mathbf{R}^d : \|x\| \leq \rho\}$$

denote the corresponding closed ball of radius ρ centered at the origin. We fix R so that $\Lambda \subset B(R)$, and we define

$$\rho_0 = \frac{R}{\|A^{-1}\|^{-1} - 1}.$$

Lemma 4.7.

- (a) Let Ω be any finite subset of Γ . Then L maps $\ell(\Omega)$ into $\ell(w_\Lambda(\Omega) \cap \Gamma)$.
- (b) If $\Omega \subset \Gamma$ is admissible, then $\ell(\Omega)$ is right-invariant under L .
- (c) $\Omega = B(\rho) \cap \Gamma$ is admissible for all $\rho \geq \rho_0$. In particular, every finite subset of Γ is contained in an admissible set.
- (d) $\Omega_\Lambda = K_\Lambda \cap \Gamma$ is admissible, and satisfies $\Omega_\Lambda = w_\Lambda(\Omega_\Lambda) \cap \Gamma$.
- (e) Let Ω be an arbitrary finite subset of Γ . If $\Omega \subset w_\Lambda(\Omega) \cap \Gamma$, then $\Omega \subset \Omega_\Lambda$.
- (f) If $\Omega \subset \Gamma$ is admissible, then $\Omega' = w_\Lambda(\Omega) \cap \Gamma \subset \Omega$ is also admissible. Further, if $\Omega_\Lambda \subset \Omega$, then $\Omega_\Lambda \subset \Omega'$.

(g) Assume $\rho \geq \rho_0$ is such that $\Omega_\Lambda \subset B(\rho)$. Then there exist admissible sets

$$\Omega_\Lambda = \Omega_0 \subsetneq \Omega_1 \subsetneq \cdots \subsetneq \Omega_{N-1} \subsetneq \Omega_N = B(\rho) \cap \Gamma$$

such that

$$w_\Lambda(\Omega_{n+1}) \cap \Gamma \subset \Omega_n, \quad n = 0, \dots, N-1. \quad (11)$$

Proof. (a) Let $a = [a_k]_{k \in \Gamma} \in \ell(\Omega)$. Since $La = [\sum_{j \in \Gamma} c_{Ai-j} a_j]_{i \in \Gamma}$ and $a_j \neq 0$ only when $j \in \Omega$, we can only have $(La)_i \neq 0$ if there is a $j \in \Omega$ such that $Ai - j \in \Lambda$. In this case, $Ai \in j + \Lambda \subset \Omega + \Lambda$, and therefore $i \in A^{-1}(\Omega + \Lambda) \cap \Gamma = w_\Lambda(\Omega) \cap \Gamma$.

(b) Follows immediately from (a).

(c) Let $\Omega = B(\rho) \cap \Gamma$. If $\rho \geq \rho_0$, then $\|A^{-1}\|(\rho + R) \leq \rho$, so

$$A^{-1}(\Omega + \Lambda) \cap \Gamma \subset A^{-1}(B(\rho + R)) \cap \Gamma \subset B(\|A^{-1}\|(\rho + R)) \cap \Gamma \subset B(\rho) \cap \Gamma = \Omega.$$

(d) Since $\Omega_\Lambda \subset K_\Lambda$, we have $w_\Lambda(\Omega_\Lambda) \cap \Gamma \subset w_\Lambda(K_\Lambda) \cap \Gamma = K_\Lambda \cap \Gamma = \Omega_\Lambda$, and therefore Ω_Λ is admissible. To show that this inclusion is an equality, suppose that $k \in \Omega_\Lambda = K_\Lambda \cap \Gamma$. By (10), there exist $\lambda_n \in \Lambda$ so that $k = \sum_{n=1}^{\infty} A^{-n} \lambda_n$. Therefore $w = Ak - \lambda_1 = \sum_{n=1}^{\infty} A^{-n} \lambda_{n+1} \in K_\Lambda$, and furthermore $w \in \Gamma$ since $Ak, \lambda_1 \in \Gamma$. Hence $w \in \Omega_\Lambda$, and therefore $k = A^{-1}(w + \lambda_1) \in A^{-1}(\Omega_\Lambda + \Lambda) \cap \Gamma = w_\Lambda(\Omega_\Lambda) \cap \Gamma$.

(e) If $\Omega \subset w_\Lambda(\Omega) \cap \Gamma = A^{-1}(\Lambda + \Omega) \cap \Gamma$, then

$$\Omega \subset A^{-1}(\Lambda) + A^{-1}(\Omega) \subset A^{-1}(\Lambda) + A^{-2}(\Lambda) + A^{-2}(\Omega) \subset \cdots$$

Since A^{-1} is a contraction, it follows that $\Omega \subset \sum_{n=1}^{\infty} A^{-n}(\Lambda) = K_\Lambda$. Since we also have $\Omega \subset \Gamma$, we conclude that $\Omega \subset \Omega_\Lambda$.

(f) If Ω is admissible, then $\Omega' = w_\Lambda(\Omega) \cap \Gamma \subset \Omega$. Hence, $w_\Lambda(\Omega') \cap \Gamma \subset w_\Lambda(\Omega) \cap \Gamma = \Omega'$, so Ω' is admissible. Further, if $\Omega_\Lambda \subset \Omega$, then $\Omega_\Lambda = w_\Lambda(\Omega_\Lambda) \cap \Gamma \subset w_\Lambda(\Omega) \cap \Gamma = \Omega'$.

(g) Define $E_0 = B(\rho)$. Since $\rho \geq \rho_0$, we have $\rho \geq \|A^{-1}\|(\rho + R)$. Therefore

$$w_\Lambda(E_0) = A^{-1}(B(\rho) + \Lambda) \subset A^{-1}(B(\rho + R)) \subset B(\|A^{-1}\|(\rho + R)) \subset B(\rho) = E_0.$$

Recursively define $E_{\nu+1} = w_\Lambda(E_\nu)$ for $\nu \geq 0$. An easy induction establishes that $E_{\nu+1} \subset E_\nu$ for every ν . Further, since E_0 is compact, we have $\cap E_\nu = K_\Lambda$ by the Contraction Mapping Theorem. Therefore, we must have $E_\nu \cap \Gamma = \Omega_\Lambda$ for all ν large enough, so $\{E_\nu \cap \Gamma\}_{\nu \geq 0}$ is a finite collection of sets. Let $\Omega_\Lambda = \Omega_0 \subsetneq \Omega_1 \subsetneq \cdots \subsetneq \Omega_N = E_0 \cap \Gamma$ be the distinct elements of this collection. Fix $0 \leq n < N$. Then there exists a ν such that $\Omega_n = E_\nu \cap \Gamma \subsetneq E_{\nu-1} \cap \Gamma = \Omega_{n+1}$. Therefore

$$w_\Lambda(\Omega_{n+1}) \cap \Gamma \subset w_\Lambda(E_{\nu-1}) \cap \Gamma = E_\nu \cap \Gamma = \Omega_n, \quad (12)$$

so (11) holds. Moreover, since $\Omega_n \subset \Omega_{n+1}$, it also follows from (12) that Ω_{n+1} is admissible. Since we know that $\Omega_0 = \Omega_\Lambda$ is admissible, the proof is complete. \square

We now show that the nonzero eigenvalues of L acting on $\ell_c(\Gamma)$ on the right coincide with the nonzero eigenvalues of L acting on $\ell(\Gamma)$ on the left, and that these further coincide with the eigenvalues of the finite submatrix

$$L_{\Omega_\Lambda} = [c_{Ai-j}]_{i,j \in \Omega_\Lambda}.$$

In particular, since L can have only finitely many nonzero eigenvalues, the accuracy of a refinable f is necessarily finite. The main ideas are adapted from one-dimensional results in [JRZ97]. The use of $\ell_c(\Gamma)$ on one side of L and $\ell(\Gamma)$ on the other is in fact natural, because $\ell_c(\Gamma)$ and $\ell(\Gamma)$ are algebraic adjoints of each other.

Let $P: \ell(\Gamma) \rightarrow (\mathbf{C}^{r \times 1})^{\Omega_\Lambda \times 1}$ denote the restriction mapping defined by

$$P([a_k]_{k \in \Gamma}) = [a_k]_{k \in \Omega_\Lambda},$$

and let $E: (\mathbf{C}^{r \times 1})^{\Omega_\Lambda \times 1} \rightarrow \ell(\Gamma)$ denote the zero extension mapping defined by

$$E([a_k]_{k \in \Omega_\Lambda}) = [a_k]_{k \in \Gamma}, \quad \text{where } a_k = 0 \text{ if } k \notin \Omega_\Lambda.$$

Theorem 4.8. Fix $\lambda \neq 0$.

- (a) If $a \in \ell_c(\Gamma)$ is a λ -eigenvector for L , then $\text{supp}(a) \subset \Omega_\Lambda$, so Pa is a λ -eigenvector for L_{Ω_Λ} . Conversely, if $a \in (\mathbf{C}^{r \times 1})^{\Omega_\Lambda \times 1}$ is a λ -eigenvector for L_{Ω_Λ} , then Ea is a λ -eigenvector for L .
- (b) If $a \in \ell(\Gamma)$ is a λ -eigenvector for L^\top , then Pa is a λ -eigenvector for $L_{\Omega_\Lambda}^\top$. Conversely, if $a \in (\mathbf{C}^{r \times 1})^{\Omega_\Lambda \times 1}$ is a λ -eigenvector for $L_{\Omega_\Lambda}^\top$, then there exists a λ -eigenvector $b \in \ell(\Gamma)$ for L^\top such that $Pb = a$.

Proof. (a) Suppose $a \in \ell_c(\Gamma)$ satisfies $La = \lambda a$ with $a \neq 0$. Set $\Omega = \text{supp}(a)$. Then $\Omega = \text{supp}(\lambda a) = \text{supp}(La) \subset w_\Lambda(\Omega) \cap \Gamma$ by Lemma 4.7(a), and therefore $\Omega \subset \Omega_\Lambda$ by Lemma 4.7(e). In particular, $Pa \neq 0$. Moreover, since $\text{supp}(a) \subset \Omega_\Lambda$, it follows that $c_{Ai-j} a_j \neq 0$ can only hold when $j \in \Omega_\Lambda$, and therefore

$$L_{\Omega_\Lambda}(Pa) = \left[\sum_{j \in \Omega_\Lambda} c_{Ai-j} a_j \right]_{i \in \Omega_\Lambda} = \left[\sum_{j \in \Gamma} c_{Ai-j} a_j \right]_{i \in \Omega_\Lambda} = P(La) = \lambda Pa.$$

Thus Pa is a λ -eigenvector for L_{Ω_Λ} .

For the converse statement, suppose that $L_{\Omega_\Lambda} a = \lambda a$ for some nonzero $a \in (\mathbf{C}^{r \times 1})^{\Omega_\Lambda \times 1}$. Note that if $j \in \Omega_\Lambda$, then $c_{Ai-j} a_j \neq 0$ can hold only if $Ai-j \in \Lambda$, which implies $i \in A^{-1}(\Omega_\Lambda + \Lambda) \cap \Gamma = \Omega_\Lambda$. Therefore,

$$L(Ea) = \left[\sum_{j \in \Omega_\Lambda} c_{Ai-j} a_j \right]_{i \in \Gamma} = E \left(\left[\sum_{j \in \Omega_\Lambda} c_{Ai-j} a_j \right]_{i \in \Omega_\Lambda} \right) = E(L_{\Omega_\Lambda} a) = \lambda Ea.$$

Therefore Ea is a λ -eigenvector for L .

(b) Choose any $a = [a_k]_{k \in \Gamma} \in \ell(\Gamma)$. Note that if $j \in \Omega_\Lambda$, then $Ai-j \in \Lambda$ can hold only if $i \in \Omega_\Lambda$ as well. Therefore,

$$\begin{aligned} PL^\top a &= P \left(\left[\sum_{i \in \Gamma} c_{Ai-j}^\top a_i \right]_{j \in \Gamma} \right) \\ &= \left[\sum_{i \in \Gamma} c_{Ai-j}^\top a_i \right]_{j \in \Omega_\Lambda} \\ &= \left[\sum_{i \in \Omega_\Lambda} c_{Ai-j}^\top a_i \right]_{j \in \Omega_\Lambda} \\ &= L_{\Omega_\Lambda}^\top Pa, \end{aligned}$$

and hence $PL^T = L_{\Omega_\Lambda}^T P$. Therefore, if $a \neq 0$ is a λ -eigenvector for L^T , then $\lambda Pa = PL^T a = L_{\Omega_\Lambda}^T Pa$. Therefore, if we can show that $Pa \neq 0$, then we can conclude that Pa is a λ -eigenvector for L^T .

Suppose that we had $Pa = 0$, i.e., that $a_k = 0$ for $k \in \Omega_\Lambda$. Since $a \neq 0$, there must be some $k_0 \in \Gamma$ such that $a_{k_0} \neq 0$. Choose $\rho \geq \rho_0$ such that $B(\rho)$ contains both Ω_Λ and k_0 , and then let $\Omega_\Lambda = \Omega_0 \subsetneq \cdots \subsetneq \Omega_N = B(\rho) \cap \Gamma$ be the admissible sets constructed in Lemma 4.7(g). Recall that Ω_N contains a point k_0 such that $a_{k_0} \neq 0$. We claim that the strictly smaller set Ω_{N-1} must also contain a point k_1 such that $a_{k_1} \neq 0$. Suppose on the contrary that $a_k = 0$ for all $k \in \Omega_{N-1}$. Note that $\lambda a_{k_0} = (L^T a)_{k_0} = \sum_{i \in \Gamma} c_{Ai-k_0}^T a_i$. However, if $Ai - k_0 \in \Lambda$, then $i \in A^{-1}(\Omega_N + \Lambda) \cap \Gamma = w_\Lambda(\Omega_N) \subset \Omega_{N-1}$, and therefore $a_i = 0$. Hence $\lambda a_{k_0} = 0$, which is a contradiction. Therefore there must be some $k_1 \in \Omega_{N-1}$ such that $a_{k_1} \neq 0$. Repeating this argument with Ω_{N-1} in place of Ω_N , we see that each Ω_n must contain a point k_n such that $a_{k_n} \neq 0$. However, since $\Omega_0 = \Omega_\Lambda$ contains no such points, this is impossible. Therefore we must have $Pa \neq 0$, so Pa is indeed a λ -eigenvector for L^T .

For the converse statement, let $a \in (\mathbf{C}^{r \times 1})^{\Omega_\Lambda \times 1}$ be a λ -eigenvector for $L_{\Omega_\Lambda}^T$. Choose $\rho \geq \rho_0$ large enough that $K_\Lambda \subset B(\rho)$, and let $\Omega_\Lambda = \Omega_0 \subsetneq \cdots \subsetneq \Omega_N = B(\rho) \cap \Gamma$ be the admissible sets constructed in Lemma 4.7(g). Set $\rho_N = \rho$, and recursively define $\rho_{n+1} = (\rho_n / \|A^{-1}\|) - R$ for $n \geq N$. If ρ is chosen large enough, then $\rho_N < \rho_{N+1} < \cdots$ is an increasing sequence of numbers. Define $\Omega_n = B(\rho_n) \cap \Gamma$ for $n > N$. By Lemma 4.7(c), these Ω_n are admissible. Combining with the sets $\Omega_0, \dots, \Omega_N$ constructed previously, we see that $\{\Omega_n\}_{n \geq 0}$ is a strictly increasing sequence of admissible sets whose union is Γ and such that $\Omega_0 = \Omega_\Lambda$. Further, by (11), we have $w_\Lambda(\Omega_{n+1}) \cap \Gamma \subset \Omega_n$ for $n = 0, \dots, N-1$. Moreover, this inclusion holds for $n \geq N$ since for these n we have

$$w_\Lambda(\Omega_{n+1}) = A^{-1}(\Omega_{n+1} + \Lambda) \subset B(\|A^{-1}\|(\rho_{n+1} + R)) = B(\rho_n).$$

Now define a_j for $j \notin \Omega_\Lambda$ recursively by the formula

$$a_j = \frac{1}{\lambda} \sum_{i \in \Omega_n} c_{Ai-j}^T a_i, \quad j \in \Omega_{n+1} \setminus \Omega_n.$$

Note that if $Ai - j \in \Lambda$ and $j \in \Omega_{n+1}$, then $i \in A^{-1}(\Omega_{n+1} + \Lambda) \cap \Gamma = w_\Lambda(\Omega_{n+1}) \subset \Omega_n$. It follows that the vector $b = [a_j]_{j \in \Gamma}$ satisfies $L^T b = \lambda b$. Since $Pb = a \neq 0$, the vector b is a λ -eigenvector for L^T . \square

We can now compare two possible finite tests for accuracy. The first test is provided by Theorem 4.4(c). To determine the accuracy corresponding to a given refinement equation, we find the largest value of p such that the finite system of linear equations in Theorem 4.4(c) has a solution. This is an iterative process: the system is solved for $s = 0, 1, \dots$ in turn until the maximum number of solvable equations is found. This test only involves solving linear equations. An upper bound on the accuracy can be found in advance by computing the right eigenvalues of $L|_{\ell(\Omega_\Lambda)}$, since accuracy p requires that $\lambda^{-\alpha}$ be an eigenvalue of L for each eigenvalue λ of A and each $0 \leq |\alpha| < p$.

An alternative test for accuracy, in the spirit of the one-dimensional results of [JRZ97], can be based on Theorem 4.4(a). Once an upper bound for p has been computed by checking the eigenvalues of $L|_{\ell(\Omega_\Lambda)}$, the left eigenvectors for L lead to the vectors $Y_{[p-1]}$ such that $Y_{[p-1]} = A_{[p-1]} Y_{[p-1]} L$. If these vectors have a polynomial structure, then the accuracy is p . If they do not have a polynomial structure, then the test must be repeated replacing p by $p-1$. This test requires the computation of the eigenvalues of a finite matrix, which cannot be done using only systems of linear equations.

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